## COMP3506/7505: Regular Exercise Set 8

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**Problem 1.** In the class, we proved that if f(h) denotes the smallest number of nodes in a balanced binary tree of height h, it must hold that

$$f(h) = 1 + f(h-1) + f(h-2).$$

Give a balanced binary tree of height 6 with f(6) nodes.

Solution.



**Problem 2.** Let T be a binary tree of n nodes. For each node u of T, define its *count* as the number of nodes in its subtree (remember that the subtree includes the node itself). Describe an algorithm to compute the counts of all the nodes in T (you can assume that each node has reserved a memory cell for you to store the count). Your algorithm must terminate in O(n) time.

**Solution.** We will design a recursive algorithm to perform a *post order* traversal of T. This algorithm at its termination will have computed the counts of all the nodes.

If the tree has only a single node, set its count to 1 and return. Otherwise, the algorithm proceeds as follows:

- Recursively compute the counts of all the nodes in the left subtree of the root.
- Recursively compute the counts of all the nodes in the right subtree of the root.
- Let  $c_1$  be the count of the left child, and  $c_2$  be the count of the right child. Set the count of the root to  $1 + c_1 + c_2$ .

To see that the running time is O(n), observe that the algorithm essentially crosses each edge of the tree twice: once from the parent to the child, and another time the other way around. There are only n-1 edges in the tree.

**Problem 3.** Let T be a binary search tree (BST) of on a set S of n integers. Let x and y be two integers in S. Describe an algorithm to find the lowest common ancestor A of the nodes in T storing x and y, respectively. If A is at level  $\ell$  (recall that the root is at level 0), your algorithm must finish in  $O(1 + \ell)$  time.

**Solution.** Without loss of generality, suppose that x < y. Let v be the root of T. If the key k of v equals either x or y, return v. Otherwise:

• If x < k < y, report v, and finish.

- If y < k, set v to its left child, and repeat the above steps at (the new) v.
- Otherwise, set v to its right child, and repeat the above steps at (the new) v.

**Problem 4.** Let T be a binary search tree (BST) of on a set S of n integers. Describe an  $O(\log n + k)$ -time algorithm to answer the following query: given an interval [a, b], report all the integers of S that fall in [a, b]. Here, k is the number of integers reported.

**Solution.** First, find the successor a' of a, and the predecessor b' of b. Then, find the lowest common ancestor, denoted as node A, of the nodes a' and b'. This takes  $O(\log n)$  time in total.

Denote by  $\Pi_1$  the path from A to node a', and  $\Pi_2$  the path from A to node b'. For every node on these two paths, report its key if the key falls in [a, b]. This takes  $O(\log n)$  time.

For every node u on  $\Pi_1$  other than A, do the following: if u is the left child of its parent p, then report all the keys in the right subtree of p. If  $k_u$  keys are reported, this step takes  $O(1 + k_u)$  time.

For every node u on  $\Pi_2$  other than A, do the following: if u is the right child of its parent p, then report all the keys in the left subtree of p. If  $k_u$  keys are reported, this step takes  $O(1 + k_u)$  time.

Overall the cost is  $O(\log n + k)$ , noticing that

$$\sum_{u\in\Pi_1\cup\Pi_2\setminus\{A\}}k_u\leq k.$$

**Problem 5.** Let S be a set of n key-value pairs of the form (t, v). Denote by m the number of distinct keys in all the pairs of S. Describe a data structure to support the following queries efficiently: given an interval [a, b], report all the pairs  $(t, v) \in S$  such that  $t \in [a, b]$ . Your structure must use O(n) space, and answer a query in  $O(\log m + k)$  time, where k is the number of pairs reported.

**Solution.** Create a BST on the m distinct keys. At each node u of the tree, use a linked list to chain up all the key-value pairs (t, v) where t equals the key of u. The query algorithm is the same as the one for Problem 4, except that for every node u whose key falls in [a, b], we should report everything in its linked list.