## COMP3506/7505: Regular Exercise Set 10

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**Problem 1.** Let T be a (2, 3)-tree on a set S of integers. Suppose that each node in the tree stores a pointer to its parent. You are given the leftmost leaf z of T, and asked to report the k smallest integers in T. Describe an algorithm to do so in O(k) time.

**Problem 2.** Let G = (V, E) be a directed graph. Suppose that we perform BFS starting from a source vertex s, and obtain a BFS-tree T. For any vertex  $v \in V$ , denote by l(v) the level of v in the BFS-tree. Prove that BFS en-queues the vertices v of V in non-descending order of l(v).

**Problem 3.** Let G = (V, E) be a directed graph. Suppose that we perform BFS starting from a source vertex s, and obtain a BFS-tree T. For any vertex  $v \in V$ , prove that the path from s to v in T is a shortest path from s to v in G.

**Problem 4.** Let G = (V, E) be an undirected graph. We will denote an edge between vertices u, v as  $\{u, v\}$ . Next, we define the single source shortest path (SSSP) problem on G. Define a path from s to t as a sequence of edges  $\{v_1, v_2\}, \{v_2, v_3\}, ..., \{v_t, v_{t+1}\}, where <math>t \ge 1, v_1 = s$ , and  $v_{t+1} = t$ . The length of the path equals t. Then, the SSSP problem gives a source vertex s, and asks to find shortest paths from s to all the other vertices in G. Adapt BFS to solve this problem in O(|V| + |E|) time. Once again, you need to produce a BFS tree where, for each vertex  $v \in V$ , the path from the root to v gives a shortest path from s to v.

**Problem 5 (Connected Components).** Let G = (V, E) be an undirected graph. A connected component (CC) of G includes a set  $S \subseteq V$  of vertices such that

- For any vertices  $u, v \in S$ , there is a path from u to v, and a path from v to u.
- (Maximality) It is not possible to add any vertex into S while still ensuring the previous property.



For example, in the above graph,  $\{a, b, c, d, e\}$  is a CC, but  $\{a, b, c, d\}$  is not, and neither is  $\{g, f, e\}$ . Prove: Let  $S_1, S_2$  be two CCs. Then, they must be disjoint, i.e.,  $S_1 \cap S_2 = \emptyset$ .

**Problem 6.** Let G = (V, E) be an undirected graph. Describe an algorithm to divide V into a set of CCs. For example, in the example of Problem 5, your algorithm should return 3 CCs:  $\{a, b, c, d, e, \}, \{g, f\}, \text{ and } \{h, i, j\}.$ 

**Problem 7.** Recall that, in the DFS algorithm, after we have grown a DFS-tree, we may need to restart from an arbitrary vertex that remains white. Multiple re-starts may be necessary throughout the algorithm. Describe how to find the re-starting white vertices efficiently so that the overall execution time of the algorithm is O(|V| + |E|).