Exercises for CSCI5010
Prepared by Yufei Tao

Problem 1. Let $P$ be a set of $n$ points in $\mathbb{R}^2$. Describe how to compute in $O(n)$ time a triangle that includes all the points of $P$ in the interior.

Hint: First compute an axis-parallel rectangle that cover all the points of $P$.

Problem 2*. Let $G = (V, E)$ be a connected regular straight-line planar graph (SLPG) with $n = |E|$ segments. Explain how to compute in $O(n)$ time a triangulated SLPG $G' = (V', E')$ such that

- $V'$ includes all the vertices of $V$, plus three dummy vertices that determine a triangle covering all the points of $V$ in interior;
- $E \subseteq E'$;
- Every face of $G'$ is covered by a face of $G$.

Hint: Identify the letmost and rightmost points in $V$. Then, find a triangle $\Delta$ that includes all the points of $V$ in the interior. The remaining obstacle is to triangulate the area “between” $G$ and the triangle’s boundary. This obstacle can be tackled by adding two segments, the first of which connects the leftmost point of $V$ to a vertex of $\Delta$, while the other connects the rightmost point of $V$ to another vertex of $\Delta$. Now, recall that an x-monotone polygon can be triangulated in linear time.

Problem 3. Let $G$ be a connected regular SLPG with $n$ segments. Describe how to build the point-location structure we discussed in $O(n \log n)$ time.

Problem 4. Prove: If a triangulated SLPG has $n$ vertices and $m$ edges, it must hold that $m = 3n - 6$.

Hint: Recall that the outer face of a triangulated SLPG is a triangle. Apply induction.

Problem 5 (Reading Exercise). Prove: The trapezoidal map defined by $n$ non-intersecting line segments in $\mathbb{R}^2$ has complexity $O(n)$.

Hint: Page 56 of Prof. Mount’s notes.

Problem 6. Describe an algorithm to build the trapezoidal map from $n$ non-intersecting line segments in $\mathbb{R}^2$ using $O(n \log n)$ time.

Problem 7*. Let $S$ be a set of $n$ non-intersecting line segments in $\mathbb{R}^2$. Given a vertical segment $q$, a query retrieves all the segments of $S$ intersecting $q$. Design a data structure of $O(n)$ space that answers a query in $O(\log n \cdot (1 + k))$ time, where $k$ is the number of segments reported. In the following example where $S = \{s_1, s_2, ..., s_5\}$, the query $q$ retrieves $k = 3$ segments.