Week 4 Tutorial

Ru Wang CSE Dept, CUHK

Adapted from the slides of the previous offerings of the course

1 of 17

Outline

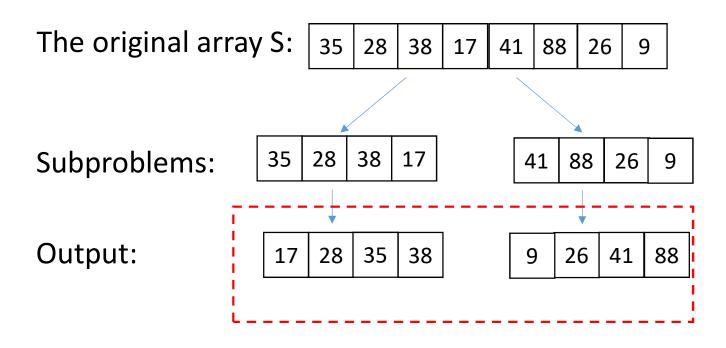
- Review recursion principle
- Review merge sort
- A variant of binary search
- A variant of merge sort
- Closest pair problem

Review – Recursion Principle

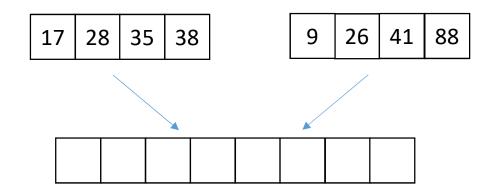
- When dealing with a subproblem (same problem but with a smaller input)
- 1. Consider it solved;
- 2. Use its output to design the rest of the algorithm.

Review – Merge Sort

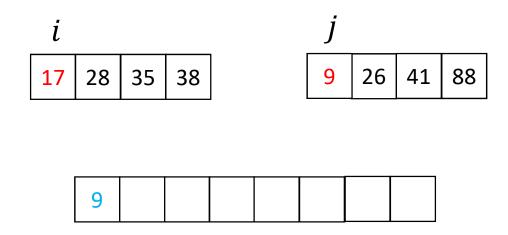
- Identify the subproblems:
 - Sort the first half of the array S.
 - Sort the second half of S.



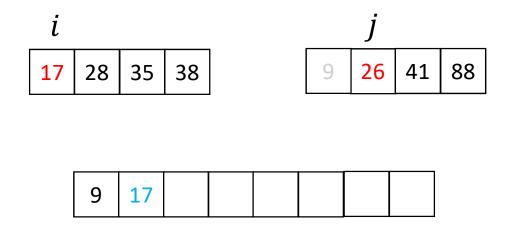
• Merge 2 sorted arrays into a single sorted array



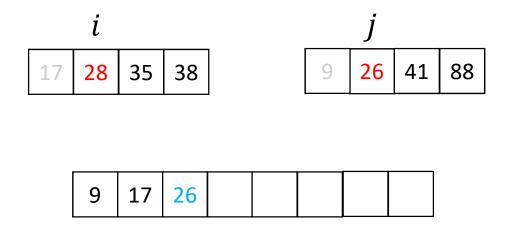
- Set *i*, *j* to 1
- Compare 17 and 9
- 9 is smaller
- Place 9 into the new array and increase *j* by 1



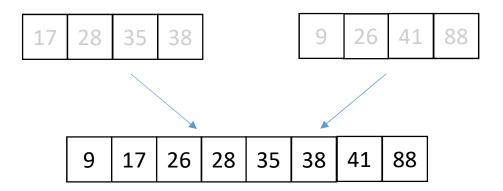
- Compare 17 and 26
- 17 is smaller
- Place 17 into the new array and increase *i* by 1



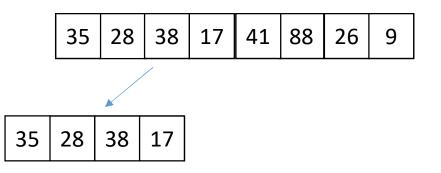
- Compare 28 and 26
- 26 is smaller
- Place 26 into the new array and increase *j* by 1

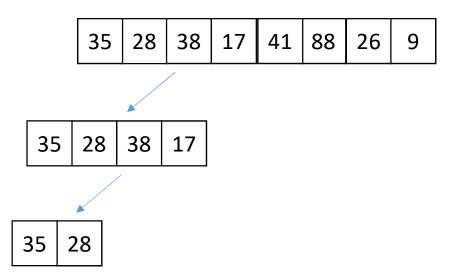


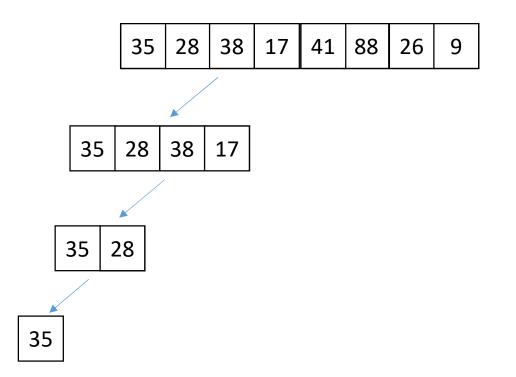
- Continue the above process until we have placed all elements into the new array
- Single pass over all the input elements
- Time complexity: O(n)

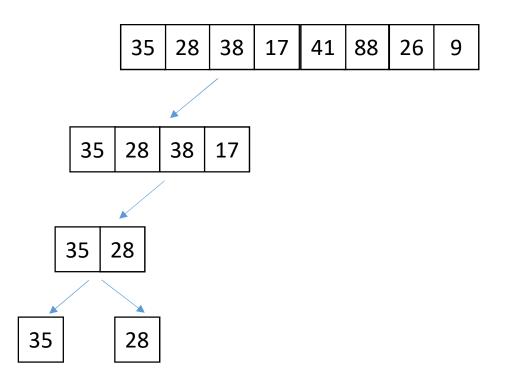


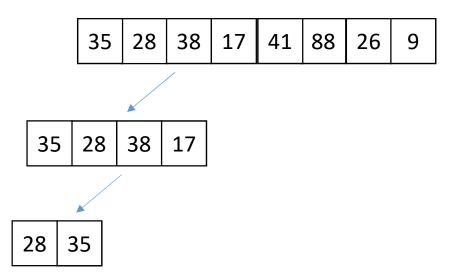
35	28	38	17	41	88	26	9
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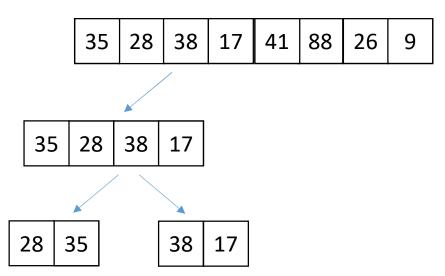


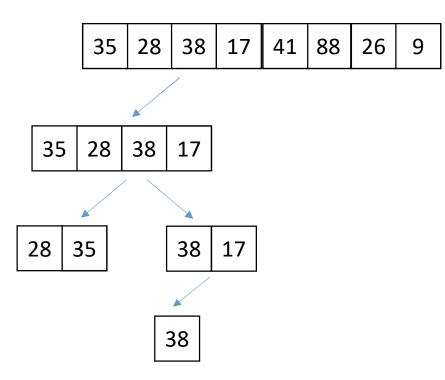


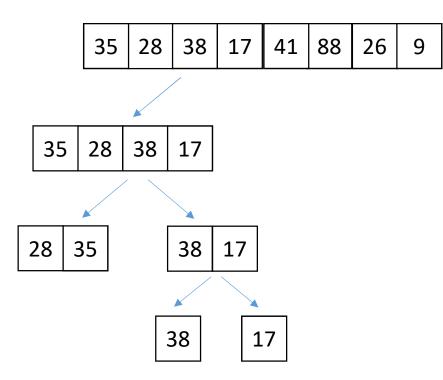


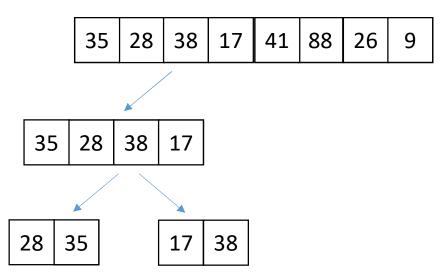


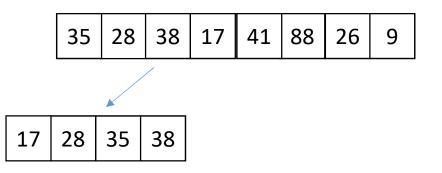


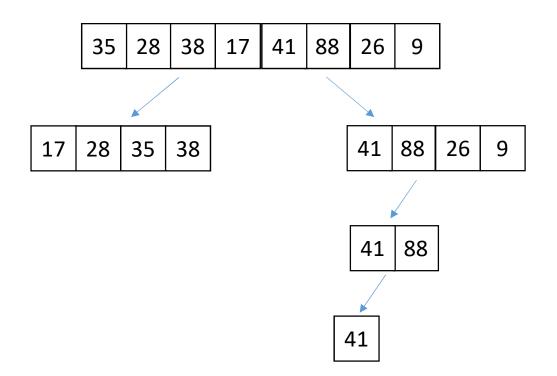


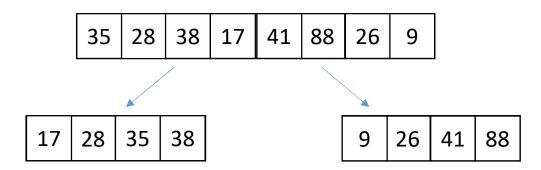








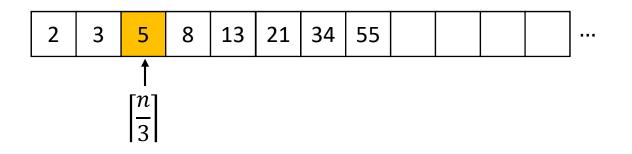




9 17 26 28	35 38	41	88	Ī
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A Variant of Binary Search

 Instead of comparing the target value with the middle element, we compare the target with the [n/3]-th element each time.



Time Complexity

- In the worst case, after each comparison, twothirds of the active elements are left.
- Solution
 - T(1) = O(1)
 - $T(n) \le T\left(\left\lceil \frac{2n}{3} \right\rceil\right) + O(1)$
 - Solving the recurrence gives $T(n) = O(\log n)$.

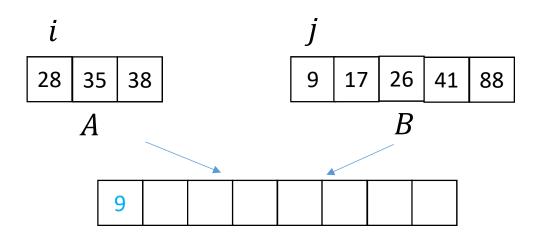
Time Complexity

- What if we compare the target with the $\left|\frac{n}{300}\right|$ -th element?
- The time complexity is still $O(\log n)$!
 - Try verifying this by yourself.
- In general, if the comparison is made to the $\left[\frac{n}{k}\right]$ -th element for some constant k > 1, the time complexity is still $O(\log n)$.

Generalized Merge Operation

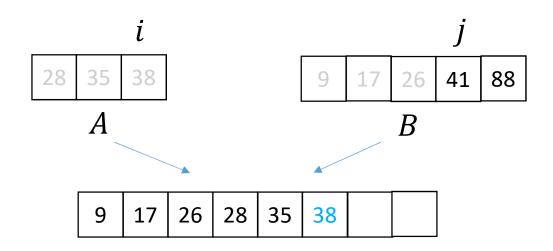
Merge 2 sorted arrays A and B, of lengths m and n.

- Set *i*, *j* to 1.
- Compare 9 with 25.
- Place 9 into the new array and increase *i* by 1.



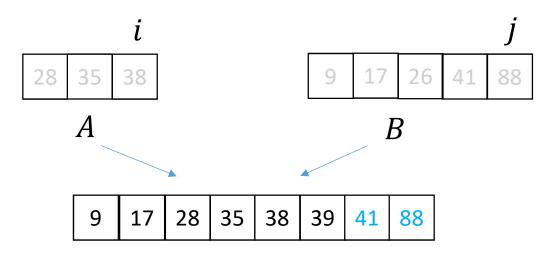
Generalized Merge Operation

• Repeat the process until we have put all the elements of one input array into the new array.



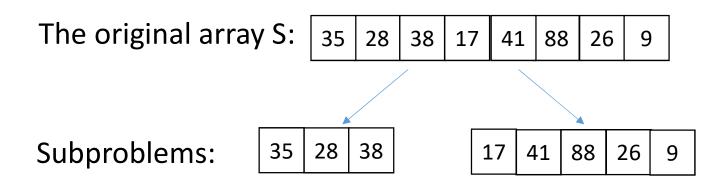
Generalized Merge Operation

- Append the remaining elements to the new array.
- Time complexity: O(m + n).



A Variant of Merge Sort

- Solve the subproblems:
 - Sort the first $\left[\frac{n}{3}\right]$ elements of the array S.
 - Sort the rest of S.
- Merge the 2 sorted arrays of different lengths.

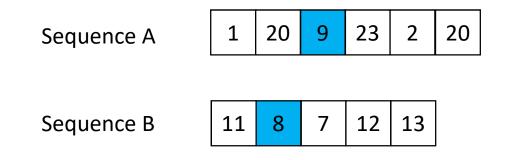


Time Complexity

- The merging takes $O\left(\left[\frac{2n}{3}\right] + \left[\frac{n}{3}\right]\right) = O(n)$ time.
- Recurrence
 - T(1) = O(1)
 - $T(n) \le T\left(\left\lceil \frac{2n}{3} \right\rceil\right) + T\left(\left\lceil \frac{n}{3} \right\rceil\right) + O(n)$
 - Solving the recurrence gives $T(n) = O(n \log n)$.
 - The recurrence can be solved with the substitution method (a regular exercise).

A Bonus Problem: Closest Pair

- Problem input:
 - Two unsorted sequences A and B with m and n integers
 - *n* < *m*
- Goal: Find a pair (x, y), x from A and y from B, with the minimum |x y|.



A Bonus Problem: Closest Pair

- This problem can be solved in $O(m \log n)$ time.
 - Sort the **shorter** sequence.
 - Then, use elements of the longer sequence to perform binary searches.
- Note: O(m log n) is better than O(m log m) when n << m.

Sequence A		20	9	23	2	20
Sorted Sequence B	7	8	11	12	13	