The RAM Computation Model

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The RAM Computation Model

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This is not a programming course.

Main take-away message from this course

Computer science is a branch of mathematics with its art reflected in the beauty of algorithms.

• Programming knowledge is not necessary to study algorithms.

Many people believe that this branch holds the future of mankind.

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In mathematics (and hence, computer science) everything — including every term and symbol — must be rigorous.

Computer science is a subject where we

- first define a computation model, which is a simple yet accurate abstraction of a computing machine;
- 2 then slowly build up a theory for this model from scratch.

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A machine has a memory and a CPU.

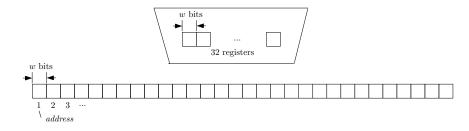
Memory

- An infinite sequence of **cells**, each of which contains the same number *w* of bits.
- Every cell has an **address**: the first cell of memory has address 1, the second cell 2, and so on.

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CPU

• Contains a fixed number — 32 in this course — of registers, each of which has w bits (i.e., same as a memory cell).



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CPU

- Can do the following atomic operations:
 - 1. (Register (Re-)Initialization)

Set a register to a fixed value (e.g., 0, -1, 100, etc.), or to the content of another register.

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CPU

• Can do the following **atomic operations**:

2. (Arithmetic)

Take the integers *a*, *b* stored in two registers, calculate one of the following, and store the result in a register:

• a + b, a - b, $a \cdot b$, and a/b.

Note: a/b is "integer division", which returns an integer. For example, 6/3 = 2 and 5/3 = 1.

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CPU

- Can do the following atomic operations:
 - 3. (Comparison/Branching)

Take the integers a, b stored in two registers, compare them, and learn which of the following is true:

• *a* < *b*, *a* = *b*, *a* > *b*.

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CPU

- Can do the following atomic operations:
 - 4. (Memory Access)

Take a memory address A currently stored in a register. Do one of the following:

- Read the content of the memory cell with address A into another register (overwriting the bits there).
- Write the content of another register into the memory cell with address *A* (overwriting the bits there).

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An **execution** is a sequence of atomic operations.

Its **cost** (also called **running time**, or simply, **time**) is the length of the sequence, namely, the number of atomic operations.

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A word is a sequence of w bits, where w is called the word length.

• In other words, each memory cell and CPU register store a word.

Unless otherwise stated, you do not need to pay attention to the value of w in this course.

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Algorithm

- An **input** refers to the initial state of the registers and the memory before an execution starts.
- An **algorithm** is a piece of description that, given an input, can be utilized to produce a sequence of atomic operations, namely, the algorithm's execution.
- The **cost** of an algorithm on an input is the length of the algorithm's execution on that input (i.e., the number of atomic operations required).
- The **space** of an algorithm on an input is the **largest** memory address accessed by the algorithm's execution on that input.

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Problem: Suppose that an integer of $n \ge 1$ has already been stored in a register. We want to calculate 1 + 2 + ... + n and store the result in a register.

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Example 1

Suppose that n is stored in register a. Set register b to 0, c to 1, and d to 1.

Repeat the following until c > a:

- Calculate b + c and store the result in b.
- Calculate c + d and store the result in c (effectively increasing c by 1).

Cost of the execution of our algorithm = 3n + 3

Think: which atomic operations are performed?

We have described the algorithm in English. The next slide shows a different way to describe the same.

Example 1

- /* *n* in register a */1. register $b \leftarrow 0, c \leftarrow 1, d \leftarrow 1$ 2. **repeat** 3. $b \leftarrow b + c$ 4. $c \leftarrow c + d$ 5. **until** c > a
- 6. return b

The above is called **pseudocode**. As you can see, we do not restrict ourselves to any particular programming languages. The description is in a "free form", mixing English words and programming-like statements as we wish, as long as it serves the purpose of clarifying —without ambiguity — our strategy.



Same problem as before. But this time we will give a faster algorithm.



Example 2

Set register *a* to 1, and *b* to 2. Let *c* be the register storing *n*. Then:

- Set a to a + c (note: now a equals n + 1).
- Set a to a * c (now a equals n(n + 1)).
- Set a to a/b (now a equals n(n+1)/2).

Cost of the execution = 5

This is significantly faster than the previous algorithm when n is large. In particular, the time of the previous algorithm increases **linearly** with n, while the time of the above one remains **constant**.

Although we have not talked about how to implement our algorithms into **actual** programs, it should be straightforward for you (who must have passed a programming course) to do so.

Computer scientists **rarely** have a programming language in mind when attacking a problem. They, instead, focus on the algorithm design because once an algorithm is clear, turning it into a program is merely a matter of translation and experience.