CSCI2100: Regular Exercise Set 3

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Problem 1. Prove $\log_2(n!) = \Theta(n \log n)$.

Problem 2. Let f(n) be a function of positive integer n. We know:

$$f(1) = 1$$

$$f(n) \le 2 + f(\lceil n/10 \rceil).$$

Prove $f(n) = O(\log n)$. Recall that $\lceil x \rceil$ is the ceiling operator that returns the smallest integer at least x.

Problem 3. Let f(n) be a function of positive integer n. We know:

$$f(1) = 1$$

$$f(n) \le 2 + f(\lceil 3n/10 \rceil).$$

Prove $f(n) = O(\log n)$. Recall that $\lceil x \rceil$ is the ceiling operator that returns the smallest integer at least x.

Problem 4. Let f(n) be a function of positive integer n. We know:

$$\begin{array}{rcl} f(1) & = & 1 \\ f(n) & \leq & 2n + 4f(\lceil n/4 \rceil). \end{array}$$

Prove $f(n) = O(n \log n)$.

Problem 5 (Bubble Sort). Let us re-visit the sorting problem. Recall that, in this problem, we are given an array A of n integers, and need to re-arrange them in ascending order. Consider the following bubble sort algorithm:

- 1. If n = 1, nothing to sort; return.
- 2. Otherwise, do the following in ascending order of $i \in [1, n-1]$: if A[i] > A[i+1], swap the integers in A[i] and A[i+1].
- 3. Recurse in the part of the array from A[1] to A[n-1].

Prove that the algorithm terminates in $O(n^2)$ time.

As an example, support that A contains the sequence of integers (10, 15, 8, 29, 13). After Step 2 has been executed once, array A becomes (10, 8, 15, 13, 29).

Problem 6* (Modified Merge Sort). Let us consider a variant of the merge sort algorithm for sorting an array A of n elements (we will use the notation A[i..j] to represent the part of the array from A[i] to A[j]):

- If n = 1 then return immediately.
- Otherwise set $k = \lceil n/3 \rceil$.
- Recursively sort A[1..k] and A[k+1..n], respectively.
- Merge A[1..k] and A[k+1..n] into one sorted array.

Prove that this algorithm runs in $O(n \log n)$ time.