Examples and Applications of Binary Search

CSCI2100 Tutorial 1

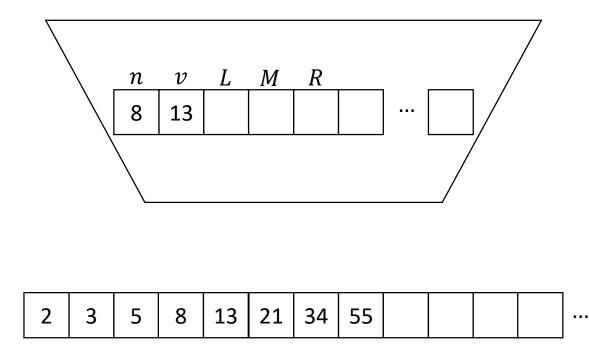
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Adapted from the slides of the previous offerings of the course

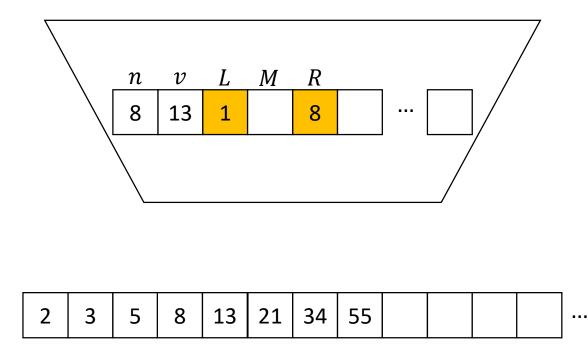
Outline

- We will first review the binary search algorithm through an example
- And then use the algorithm to solve a "two-sum" problem.

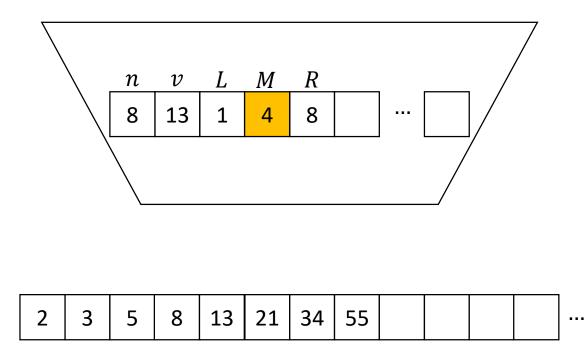
• Suppose we have the following sorted input set S, and are trying to find the value 13.



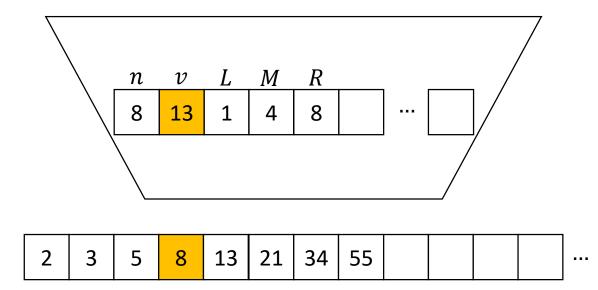
• Initializing L to be 1 and R to n (in this case, 8)



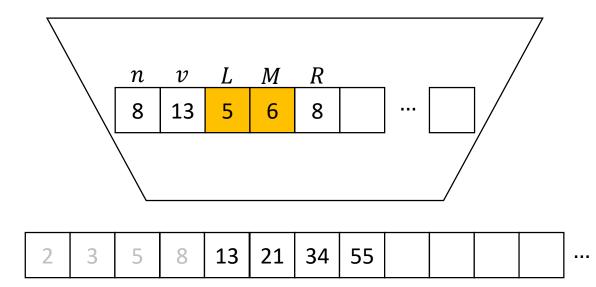
- Since $L \leq R$
- Proceed by computing M = (L + R)/2



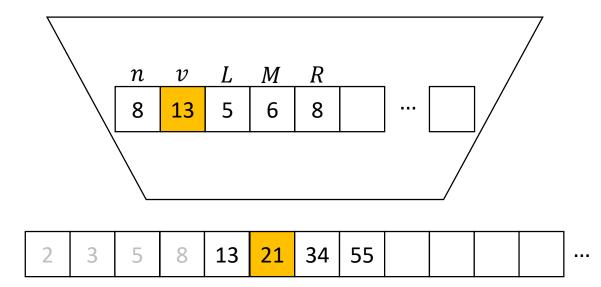
- Compare v = 13 and the value 8 indexed by M
- v > the value indexed by M
- Means that the target is in the right half of the sorted sequence



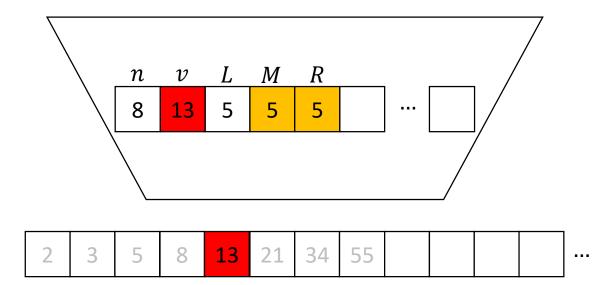
- Look at the right half of the sorted sequence
- Set L to be M + 1 (discard the left half)
- Recompute *M*



- Compare v and the value 21 indexed by M
- v < the value indexed by M
- Means that the target is in the left half of the sorted sequence



- Set R to be M 1 (discard the right half)
- *L*, *R*, M = 5
- v = the value indexed by M, return "yes"



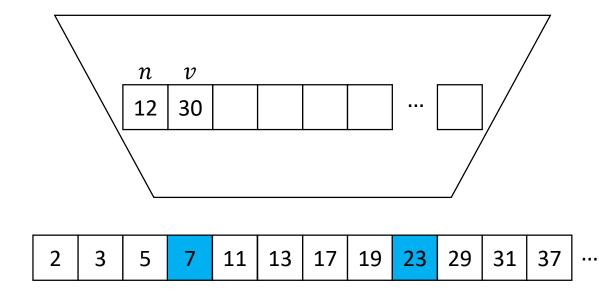
The Two-Sum Problem

- Problem Input:
 - A sequence of *n* positive integers in strictly increasing order in memory at the cells numbered from 1 up to *n*
 - The value *n* has been placed in Register 1
 - A positive integer v has been placed in Register 2
- Goal:
 - Determine whether if there exist two different integers x and y in the sorted sequence such that x + y = v

2	3	5	7	11	13	17	19	23	29	31	37	
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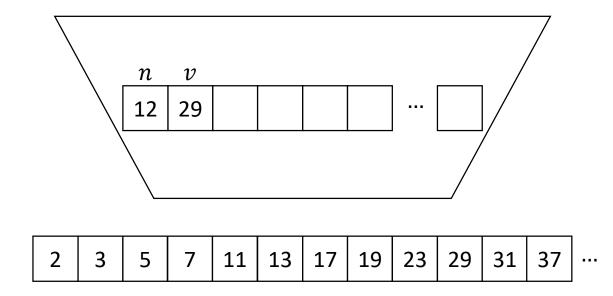
Example

• A "yes"-input with n = 12, v = 30



Example

• A "no"-input with n = 12, v = 29



A First Attempt

- Naïve algorithm:
 - Enumerate all possible pairs in the sorted sequence
 - Check if they sum to v
 - There are $\binom{n}{2} = \frac{n(n-1)}{2}$ possible pairs
 - Worst-case time: at least n(n-1)/2
- Can we do better than this?
- Hint: Take advantage of the fact that the given sequence is sorted!

Binary Search the Answer

- Goal: Find a pair(x, y) such that x + y = v
- Observe that given x, y = v x, is determined
- Improve the naïve algorithm
 - Instead of enumerating all possible y , we can find if there exits an integer $\boldsymbol{v}-\boldsymbol{x}$ in the sequence
- Solution:
 - For each *x* in the sequence:
 - set y as v x
 - Use binary search to see if y exists in the sequence

The Repeated Binary Search Algorithm

- Pseudocode:
 - 1. Let *n* be register 1 and *v* be register 2
 - 2. register $i \leftarrow 1$, register $one \leftarrow 1$
 - 3. while $i \leq n$
 - 4. read into register x the memory cell at address *i*

5.
$$y \leftarrow v - x$$

6. **if**
$$BinarySearch(y) = "yes"$$

- 7. return "yes"
- 7. $i \leftarrow i + one$ (effectively increasing *i* by 1)
- 8. return "no"

Worst-Case Running Time

- Worst case (when the output is "no")
- This algorithm needs to run binary search *n* times
- Cost of each binary search: at most $10(1 + \log_2 n)$
- Cost of the algorithm: at most $100n(1 + \log_2 n)$ (a loose upper bound)
- Can we do even better?
- Actually this problem can be solved in at most 100*n* time --- left for you to try outside the class.