# Examples and Applications of Binary Search <br> CSCI2100 Tutorial 1 <br> Shangqi Lu 

Adapted from the slides of the previous offerings of the course

## Outline

- We will first review the binary search algorithm through an example
- And then use the algorithm to solve a "two-sum" problem.


## Binary Search Review

- Suppose we have the following sorted input set S , and are trying to find the value 13.



## Binary Search Review

- Initializing $L$ to be 1 and $R$ to $n$ (in this case, 8 )



## Binary Search Review

- Since $L \leq R$
- Proceed by computing $M=(L+R) / 2$



## Binary Search Review

- Compare $v=13$ and the value 8 indexed by $M$
- $v>$ the value indexed by $M$
- Means that the target is in the right half of the sorted sequence


| 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\ldots$ |  |  |  |  |  |  |  |  |  |  |  |

## Binary Search Review

- Look at the right half of the sorted sequence
- Set $L$ to be $M+1$ (discard the left half)
- Recompute $M$


| 2 | 3 | 5 | 8 | 13 | 21 | 34 | 55 |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\ldots$ |  |  |  |  |  |  |  |  |  |  |  |

## Binary Search Review

- Compare $v$ and the value 21 indexed by $M$
- $v<$ the value indexed by $M$
- Means that the target is in the left half of the sorted sequence



## Binary Search Review

- Set $R$ to be $M-1$ (discard the right half)
- $L, R, M=5$
- $v=$ the value indexed by $M$, return "yes"



## The Two-Sum Problem

- Problem Input:
- A sequence of $n$ positive integers in strictly increasing order in memory at the cells numbered from 1 up to $n$
- The value $n$ has been placed in Register 1
- A positive integer $v$ has been placed in Register 2
- Goal:
- Determine whether if there exist two different integers $x$ and $y$ in the sorted sequence such that $x+y=v$

| 2 | 3 | 5 | 7 | 11 | 13 | 17 | 19 | 23 | 29 | 31 | 37 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\ldots$ |  |  |  |  |  |  |  |  |  |  |  |

## Example

- A "yes"-input with $n=12, v=30$



## Example

- A "no"-input with $n=12, v=29$



## A First Attempt

- Naïve algorithm:
- Enumerate all possible pairs in the sorted sequence
- Check if they sum to $v$
- There are $\binom{n}{2}=\frac{n(n-1)}{2}$ possible pairs
- Worst-case time: at least $n(n-1) / 2$
- Can we do better than this?
- Hint: Take advantage of the fact that the given sequence is sorted!


## Binary Search the Answer

- Goal: Find a pair $(x, y)$ such that $x+y=v$
- Observe that given $\mathrm{x}, y=v-x$, is determined
- Improve the naïve algorithm
- Instead of enumerating all possible $y$, we can find if there exits an integer $v-x$ in the sequence
- Solution:
- For each $x$ in the sequence:
- set $y$ as $v-x$
- Use binary search to see if $y$ exists in the sequence


## The Repeated Binary Search Algorithm

- Pseudocode:

1. Let $n$ be register 1 and $v$ be register 2
2. register $i \leftarrow 1$, register one $\leftarrow 1$
3. while $i \leq n$
4. read into register $x$ the memory cell at address $i$
5. $\quad y \leftarrow v-x$
6. if BinarySearch $(y)=$ "yes"
7. return "yes"
8. $\quad i \leftarrow i+$ one (effectively increasing $i$ by 1 )
9. return "no"

## Worst-Case Running Time

- Worst case (when the output is "no")
- This algorithm needs to run binary search $n$ times
- Cost of each binary search: at most $10\left(1+\log _{2} n\right)$
- Cost of the algorithm: at most $100 n\left(1+\log _{2} n\right)$ (a loose upper bound)
- Can we do even better?
- Actually this problem can be solved in at most $100 n$ time --- left for you to try outside the class.

