Basic Concepts and Representation Methods of Graphs

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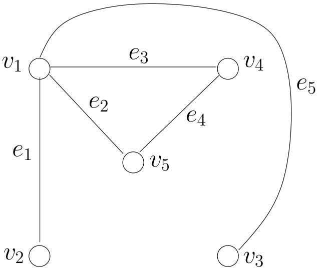
Undirected Graphs

An undirected graph is a pair of (V, E) where:

- V is a set of elements, each of which called a node.
- E is a set of **unordered pairs** $\{u, v\}$ such that u and v are nodes.

A node may also be called a vertex. We will refer to V as the vertex set or the node set of the graph, and E the edge set.

Example



This is an undirected graph where there are 5 vertices $v_1, v_2, ..., v_5$, and 5 edges $e_1, e_2, ..., e_5$.

Directed Graphs

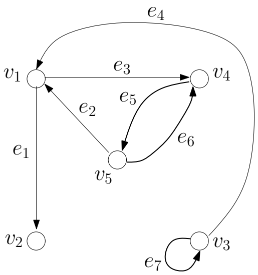
An directed graph is a pair of (V, E) where:

- V is a set of elements, each of which called a node.
- E is a set of pairs (u, v) where u and v are nodes in V. We say that there is a (directed) edge from u to v.

A node may also be called a vertex. We will refer to V as the vertex set or the node set of the graph, and E the edge set.

A (directed) edge (u, v) is said to be an outgoing edge of u, and an incoming edge of v. Accordingly, v is an out-neighbor of u, and u an in-neighbor of v.

Example

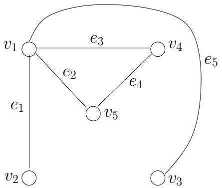


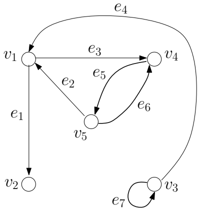
This is an directed graph (V, E) where there are 5 vertices $v_1, v_2, ..., v_5$, and 7 edges $e_1, e_2, ..., e_7$. Note that every edge has a direction. Edge e_6, for instance, is an outgoing edge of v_5, and an incoming edge of v_4.

Degrees

- In an undirected graph, the degree of a vertex u is the number of edges of u.
- In a directed graph, the out-degree of a vertex u is the number outgoing edges of u, and its in-degree is the number of its incoming edges.

Example





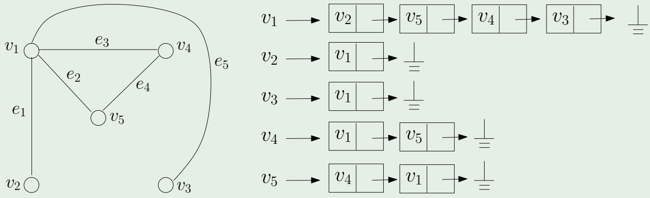
In the left graph, the degree of v_5 is 2. In the right graph, the out-degree of v_3 is 2, and its in-degree is 1.

Next, we discuss two common ways to store a graph: **adjacency list** and **adjacency matrix**. In both cases, we represent each vertex in V using a unique id in $1, 2, ..., |V|$.

Adjacency List – Undirected Graphs

Each vertex $u \in V$ is associated with a linked list that enumerates all the vertices that are connected to u.

Example 1.

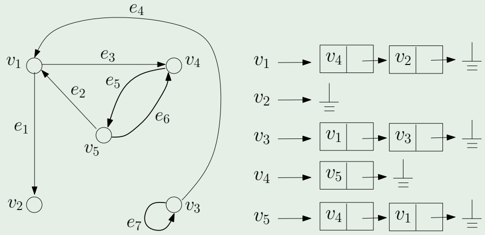


Space $= O(|V| + |E|)$.

Adjacency List – Directed Graphs

Each vertex $u \in V$ is associated with a linked list that enumerates all the vertices $v \in V$ such that there is an edge from u to v.

Example 2.

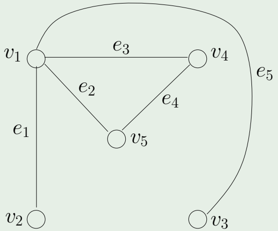


Space $= O(|V| + |E|)$.

Adjacency Matrix – Undirected Graphs

A $|V| \times |V|$ matrix A where $A[u, v] = 1$ if $(u, v) \in E$, or 0 otherwise.

Example 3.



	v_1	v_2	v_3	v_4	v_5
v_1	0	1	1	1	1
v_2	1	0	0	0	0
v_3	1	0	0	0	0
v_4	1	0	0	0	1
v_5	1	0	0	1	0

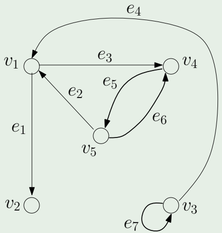
- A must be symmetric.
- Space $= O(|V|^2)$.

Think: How to store A so that, for any vertices $u, v \in V$, we can find out if they have an edge in constant time?

Adjacency Matrix – Directed Graphs

Defined in the same way as in the undirected case.

Example 4.



	v_1	v_2	v_3	v_4	v_5
v_1	0	1	0	1	0
v_2	0	0	0	0	0
v_3	1	0	1	0	0
v_4	0	0	0	0	1
v_5	1	0	0	1	0

- A may not be symmetric.
- Space $= O(|V|^2)$.