CSCI2100: Regular Exercise Set 4

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Problem 1. Recall that our RAM model has been extended with an atomic operation RANDOM(x, y) which, given integers x, y, returns an integer chosen uniformly at random from [x, y]. Suppose that you are allowed to call the operation *only* with x = 1 and y = 128. Describe an algorithm to obtain a uniformly random number between 1 and 100. Your algorithm must finish in O(1) expected time.

Problem 2*. Suppose that we enforce an even harder constraint that you are allowed to call RANDOM(x, y) only with x = 0 and y = 1. Describe an algorithm to generate a uniformly random number in [1, n] for an arbitrary integer n. Your algorithm must finish in $O(\log n)$ expected time.

Problem 3. For the k-selection problem, consider an input array A that has n = 120 elements. Our randomized algorithm selects a number v, and recurse into a smaller array A' if the rank of v is within [n/3, 2n/3] = [40, 80]. For k = 20, what is the probability that the size of A' is at most 60?

Problem 4* (Textbook Exercise 9.3-8). Let X[1..n] and Y[1..n] be two arrays, each containing n integers in ascending order. Consider that all the 2n integers are distinct. Let k be an integer between 1 and 2n. Give an $O(\log n)$ -time algorithm for finding the k-th smallest of the 2n elements.

Problem 5^{**} (A Simpler Randomized Algorithm for k-Selection, but with a More Tedious Analysis). In the k-selection problem, we have an array S of n distinct integers (not necessarily sorted). We would like to find the k-th smallest integer in S where $k \in [1, n]$. Here is another way of solving it using randomization. If n = 1, then we simply return the only element in S. For n > 1, we proceed as follows:

- Randomly pick an integer v in S, and obtain the rank r of v in S.
- If r = k, return v.
- If r > k, produce an array S' containing the integers of S that are smaller than v. Recurse by finding the k-th smallest in S'.
- Otherwise, produce an array S' containing the integers of S that are larger than v. Recurse by finding the (r k)-th smallest in S'.

Prove that the above algorithm finishes in O(n) expected time.