## CSCI2100: Regular Exercise Set 4

Prepared by Yufei Tao

Problem 1. Recall that our RAM model has been extended with an atomic operation RANDOM $(x, y)$ which, given integers $x, y$, returns an integer chosen uniformly at random from $[x, y]$. Suppose that you are allowed to call the operation only with $x=1$ and $y=128$. Describe an algorithm to obtain a uniformly random number between 1 and 100 . Your algorithm must finish in $O(1)$ expected time.

Problem 2*. Suppose that we enforce an even harder constraint that you are allowed to call $\operatorname{RANDOM}(x, y)$ only with $x=0$ and $y=1$. Describe an algorithm to generate a uniformly random number in $[1, n]$ for an arbitrary integer $n$. Your algorithm must finish in $O(\log n)$ expected time.

Problem 3. For the $k$-selection problem, consider an input array $A$ that has $n=120$ elements. Our randomized algorithm selects a number $v$, and recurse into a smaller array $A^{\prime}$ if the rank of $v$ is within $[n / 3,2 n / 3]=[40,80]$. For $k=20$, what is the probability that the size of $A^{\prime}$ is at most 60 ?

Problem 4* (Textbook Exercise 9.3-8). Let $X[1 . . n]$ and $Y[1 . . n]$ be two arrays, each containing $n$ integers in ascending order. Consider that all the $2 n$ integers are distinct. Let $k$ be an integer between 1 and $2 n$. Give an $O(\log n)$-time algorithm for finding the $k$-th smallest of the $2 n$ elements.

Problem 5** (A Simpler Randomized Algorithm for k-Selection, but with a More Tedious Analysis ). In the $k$-selection problem, we have an array $S$ of $n$ distinct integers (not necessarily sorted). We would like to find the $k$-th smallest integer in $S$ where $k \in[1, n]$. Here is another way of solving it using randomization. If $n=1$, then we simply return the only element in $S$. For $n>1$, we proceed as follows:

- Randomly pick an integer $v$ in $S$, and obtain the rank $r$ of $v$ in $S$.
- If $r=k$, return $v$.
- If $r>k$, produce an array $S^{\prime}$ containing the integers of $S$ that are smaller than $v$. Recurse by finding the $k$-th smallest in $S^{\prime}$.
- Otherwise, produce an array $S^{\prime}$ containing the integers of $S$ that are larger than $v$. Recurse by finding the $(r-k)$-th smallest in $S^{\prime}$.

Prove that the above algorithm finishes in $O(n)$ expected time.

