

## Exercises for CSCI5010

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**Problem 1.** Let  $P$  be a set of  $n$  points in  $\mathbb{R}^d$ . Define  $d_{max}$  as the maximum distance between two points in  $P$ , namely,  $d_{max} = \max_{p,q \in P} \text{dist}(p,q)$ , where  $\text{dist}(p,q)$  is the Euclidean distance between  $p$  and  $q$ . Describe an  $O(n)$ -time algorithm to find a  $d$ -dimensional box such that:

- the box has the same side length on each dimension, and the side length is  $\Theta(d_{max})$ ;
- the box covers all the points in  $P$ .

**Problem 2.** Let  $P$  be a set of  $n$  points in  $\mathbb{R}^d$ . Define  $d_{max}$  as the maximum distance between two points in  $P$ , namely,  $d_{max} = \max_{p,q \in P} \text{dist}(p,q)$ . Define  $d_{min}$  as the minimum distance between two distinct points in  $P$ , namely,  $d_{min} = \min_{\text{distinct } p,q \in P} \text{dist}(p,q)$ . Describe how to build a quadtree on  $P$  in  $O(n \cdot \log \frac{d_{max}}{d_{min}})$  time.

**Problem 3.** Let  $P$  be a set of points in  $\mathbb{R}^d$ . Suppose that we have constructed an  $s$ -well separated point decomposition for  $P$ :  $\{\{A_1, B_1\}, \{A_2, B_2\}, \dots, \{A_m, B_m\}\}$ . Let  $\{p, q\}$  be a closest pair of  $P$  (i.e.,  $\text{dist}(p, q) = d_{min}$ , where  $d_{min}$  is as defined in Problem 2). Prove: if  $s > 2$ , then there exists an  $i \in [1, m]$  such that  $A_i$  contains only  $p$ , and  $B_i$  contains only  $q$ .

Hint 1: Recall that there must exist an  $i \in [1, m]$  such that  $p \in A_i$  and  $q \in B_i$ .

Hint 2: If you do not want to think, read Lect 18 of Prof. Mount's notes.

**Problem 4\*.** Let  $P$  be a set of points in  $\mathbb{R}^d$ . Suppose that we have constructed an  $s$ -well separated point decomposition for  $P$ :  $\{\{A_1, B_1\}, \{A_2, B_2\}, \dots, \{A_m, B_m\}\}$ . For each  $i \in [1, m]$ , let  $a_i$  be an arbitrary point in  $A_i$ , and  $b_i$  be an arbitrary point in  $B_i$ . Let us construct an undirected graph  $G = (V, E)$  as follows:

- $V = P$ , namely, each vertex of  $V$  is a point in  $P$ .
- For each  $i \in [1, m]$ , add to  $E$  an edge  $\{a_i, b_i\}$ .

Prove: If  $s > 2$ , then  $G$  must be connected (i.e., for any two points  $p, q \in P$ ,  $G$  has a path from  $p$  to  $q$ ).

Hint: Imagine listing all distinct pairs of the points in  $P$  in ascending order of distance. Apply induction on the sorted list.