## CSCI5010: Final Exam

Name:

## Student ID:

Please write all your solutions in the answer book, except Problems 1 and 2 which should be answered in this paper.

Problem 1 (10\%). The figure below shows a set of 5 segments. Give the trapezoidal map that is decided by these segments.


Problem 2 ( $\mathbf{1 0 \%}$ ). The left figure below shows the Delaunay triangulation of the set of black points. Suppose that we want to insert point $p$ (i.e., the white point). Draw the resulting Delaunay triangulation in the figure on the right.


Problem 3 (20\%). Let $P$ be a set of $n$ points in $\mathbb{R}^{2}$. Given an axis-parallel rectangle $q$, a query reports the number of points in $q \cap P$. Describe a data structure of $O(n)$ size that answers such a query in $O(\sqrt{n})$ time.

Problem $4 \mathbf{( 2 0 \% )}$. Let $S$ be a set of horizontal segments in $\mathbb{R}^{2}$, where each segment has the form $\left[x_{1}, x_{2}\right] \times y$. Given a point $q$, a query reports the first segment of $S$ that will be hit if we shoot a ray
upwards from $q$ (e.g., in the figure below, the query reports $s$ ). Preprocess $S$ into a data structure of $O(n)$ space such that a query can be answered in $O(\log n)$ time.


Problem 5 (20\%). Let $S$ and $T$ be two sets of points in $\mathbb{R}^{2}$. Let $(p, q)$ be a closest pair of $S$ and $T$, namely, the Euclidean distance between $p$ and $q$ is the smallest among all pairs of points in $S \times T$. For example, in the figure below, let $S(T)$ be the set of black (white) points. The closest pair is the two points between which there is a segment. Prove that there must be an edge between $p$ and $q$ in the Delaunay triangulation of $S \cup T$.


Problem 6 (20\%). Let $P$ be a set of $n$ points in $\mathbb{R}^{2}$. Given a rectangle $r$ and a query point $q$, a constrained nearest neighbor query returns the point in $P \cap r$ that has the smallest Euclidean distance to $q$ (i.e., among all the points of $P$ falling in $r$, report the one closest to $q$ ). For example, in the figure below, let $P$ be the set of black points; given the rectangle $r$ and $q$ as shown, a query returns point $p_{1}$ as its answer (note that the answer is not $p_{2}$ as it is outside $r$ ). Give a structure of $O\left(n \log ^{2} n\right)$ space that answers such a query in $O\left(\log ^{3} n\right)$ time.


