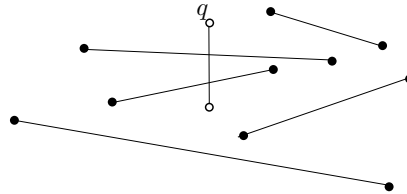


## CSCI5010 Exercise List 10

**Problem 1 (Segment-Segment Intersection).** Let  $S$  be a set of non-intersecting segments in  $\mathbb{R}^2$ . Given a vertical segment  $q = x \times [y_1, y_2]$ , a query reports all the segments in  $S$  intersecting  $q$  (e.g., in the figure below, the query reports two segments). Preprocess  $S$  into a data structure of  $O(n)$  space such that a query can be answered in  $O((k + 1) \log n)$  time, where  $k$  is the number of segments reported.



**Problem 2 (Segment-Segment Intersection).** Consider the previous problem again. Let  $s, s'$  be two segments in  $S$ . We say that  $s'$  is *vertically adjacent* to  $s$  if we can shoot a vertical ray either upward or downward from an endpoint of  $s'$ , such that  $s$  is the first segment in  $S$  hit by the ray. The *vertical visibility* of  $s$  is defined as the number of segments in  $S$  (other than  $s$ ) vertically adjacent to  $s$ . For example, the vertical visibility of  $s_1$  is 3, while that of  $s_2$  is 1.

Let  $\lambda$  be the maximum vertical visibility of all segments in  $S$ . Preprocess  $S$  into a data structure of  $O(n)$  space such that a query can be answered in  $O(\log n + k \log \lambda)$  time, where  $k$  is the number of segments reported.

