## CSCI3610: Special Exercise Set 1

**Problem 1.** Explain how to implement the operation  $x \mod y$  in O(1) time where x and y are positive integers.

**Problem 2.** For the k-selection problem, suppose that the input is an array of 12 elements: (58, 23, 98, 83, 32, 24, 18, 45, 85, 91, 2, 34). Recall that our randomized algorithm first selects a number v and then recursively solves a subproblem. Suppose that v = 34 and k = 10. What is the size of the array for the subproblem?

**Problem 3 (Textbook Exercise 9.3-5).** The *median* of a set S of n elements is the  $\lfloor n/2 \rfloor$  smallest element in S. Suppose that you are given a deterministic algorithm for finding the median of S (stored in an array) in O(n) worst-case time. Using this algorithm as a black box, design another deterministic algorithm for solving the k-selection problem (for any  $k \in [1, n]$ ) in O(n) worst-case time.

**Problem 4.** Let S be a set of n distinct integers, and  $k_1, k_2$  be arbitrary integers satisfying  $1 \le k_1 \le k_2 \le n$ . Suppose that S is given in an array. Give an O(n) expected time algorithm to report *all* the integers whose ranks in S are in the range  $[k_1, k_2]$ . Recall that the rank of an integer v in S equals the number of integers in S that are at most v.

**Problem 5\* (Textbook Exercise 9-2).** We are given an array that stores a set S of n distinct positive integers. Set  $W = \sum_{e \in S} e$ . Describe an algorithm to find the element  $e^* \in S$  that makes both of the following hold:

- $\sum_{e < e^*} e < W/2$
- $\sum_{e>e^*} e \leq W/2$ .

Your algorithm should finish in O(n) time (O(n) expected time is acceptable).

(Hint: First convince yourself that such  $e^*$  is unique, and then adapt the k-selection algorithm).