

## CSCI2100: Midterm Solutions

**Problem 1 (10%).**  $5n + 3\sqrt{n} \leq 8n$  for all  $n \geq 1$ .

**Problem 2 (12%).** Use the  $k$ -selection algorithm to find the  $k_1$ -th smallest integer  $a_1$  of  $S$ . Then, run the algorithm again to find the  $k_2$ -th smallest integer  $a_2$  of  $S$ . Finally, scan  $S$  to report all the integers  $x \in S$  satisfying  $a_1 \leq x \leq a_2$ .

**Problem 3 (13%).** Create a hash table on  $S_1$  in  $O(n)$  time. Then, for each  $x \in S_2$ , probe the hash table to see whether  $x \in S_1$ ; if so, output it.

**Problem 4 (10%).** First, it is clear that 10 is the root. As  $10 < 30$ , we must be entering the right subtree of 10. As  $30 > 20$ , we must be entering the left subtree of 30. As the next number is 60, we know that 60 lies in the left subtree of 30. This is not possible in a BST.

**Problem 5 (10%).** The statement is correct. Let  $u$  (resp.,  $v$ ) be the left (resp., right) child node of the root. Denote by  $k$  the 2nd smallest integer of  $S$ . If  $k$  is at neither  $u$  nor  $v$ , then it must be the key of a proper descendant of either  $u$  or  $v$ . In the former case,  $k$  is smaller than the key of  $u$ , which is not possible in a heap. Similarly, the latter case is not possible, either.

**Problem 6 (15%).** Initialize an empty linked list  $Q$ . At all times,  $Q$  stores its numbers in ascending order. To perform an `ins-large`( $e$ ), we check if  $e$  is larger than the integer at the tail of  $Q$ . If not, ignore  $e$ ; otherwise, append  $e$  to the end of  $Q$ . To perform a `del-min`, simply remove the first element of  $Q$ .

**Problem 7 (15%).** We first solve the following “two-sum” problem (in fact, the problem was discussed in a tutorial). Given an integer  $t$ , determine if there exist distinct  $i, j \in [1, n]$  such that  $A[i] + A[j] = t$ . For this purpose, first create a hash table on the elements in  $A$ , which takes  $O(n)$  time. Then, for each  $i \in [1, n]$ , check whether  $t - A[i]$  exists in the hash table, which takes  $O(1)$  expected time. Doing so for all  $i \in [1, n]$  requires  $O(n)$  expected time.

To solve the original problem, we will deal with  $n$  instances of the two-sum problem. Specifically, for each  $i \in [1, n]$ , check whether there exist distinct  $j, k \in [1, n] \setminus \{i\}$  such that  $A[j] + A[k] = t - A[i]$ . As discussed, each instance incurs  $O(n)$  expected time. The total amount of time is thus  $O(n^2)$  in expectation.

**Problem 8 (15%).** Initialize an empty min-heap  $H'$ . Insert the key of the root of  $H$  into  $H'$ . Repeat the following procedure  $k$  times.

- Perform a `del-min` from  $H'$ . Denote by  $u$  the node removed by the `del-min` operation. Report the key of  $u$ .
- Let  $v_1$  and  $v_2$  be the child nodes of  $u$  in  $H$ . Insert  $v_1$  and  $v_2$  to  $H'$ .

The heap  $H'$  has at most  $k$  elements (every `del-min` on  $H'$  is followed by at most two insertions into  $H'$ ). Therefore, every execution of the above procedure takes  $O(\log k)$  time, resulting in  $O(k \log k)$  time overall.