Quick Sort

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Today, we will discuss another sorting algorithm named **quick sort**. It is a randomized algorithm that runs in $O(n^2)$ time in the **worst** case but $O(n \log n)$ time **in expectation**.

Recall:

The Sorting Problem

Problem Input:

A set S of n integers is given in an array A of length n.

Goal:

Produce an array that stores the elements of S in ascending order.

Quick Sort

1. Pick an integer p from A **uniformly at random**, which is called the **pivot**.
2. Store the integers in another array A' such that
   - all the integers smaller than p are before p in A';
   - all the integers larger than p are after p in A'.
3. Sort the part of A' before p recursively (a subproblem).
4. Sort the part of A' after p recursively (a subproblem).

Example

After Step 1 (suppose that 26 was randomly picked as the pivot):

				p																									
38	28	88	17	26	41	72	83	69	47	12	68	5	52	35	9														

After Step 2:

				p																									
17	12	5	9	26	38	28	88	41	72	83	69	47	68	52	35														

After Steps 3 and 4:

				p																									
5	9	12	17	26	28	35	38	41	47	52	68	69	72	83	88														

Analysis of Quick Sort

Quick sort is not attractive in the worst case: its worst case time is $O(n^2)$ (why?). However, quick sort is fast in expectation: we will prove that its expected time is $O(n \log n)$. Remember: this holds on **every** input array A.

The rest of the slides will not be tested for CSCI2100.

Analysis of Quick Sort

First, convince yourself that it suffices to analyze the number X of comparisons. The running time is bounded by $O(n + X)$.

Next, we will prove that $\boldsymbol{E}[X] = O(n \log n)$.

Analysis of Quick Sort

Denote by e_i the i-th smallest integer in S. Consider e_i, e_j for any i, j such that $i \neq j$.

What is the probability that quick sort compares e_i and e_j?

This question, which seems to be difficult at first glance, has a surprisingly simple answer. Let us observe:

- Every element will be selected as a pivot **exactly** once.
- e_i and e_j are **not** compared, if any element **between** them gets selected as a pivot **before** e_i and e_j.

> For example, suppose that $i = 7$ and $j = 12$. If e_9 is the pivot, then e_i and e_j will be separated by e_9 (**think:** why?) and will not be compared in the rest of the algorithm.

Analysis of Quick Sort

Therefore, e_i and e_j are compared if and only if either one is the first among $e_i, e_{i+1}, ..., e_j$ picked as a pivot.

The probability is $2/(j - i + 1)$.

Analysis of Quick Sort

Define random variable X_{ij} to be 1, if e_i and e_j are compared. Otherwise, $X_{ij} = 0$. We thus have $\boldsymbol{Pr}[X_{ij} = 1] = 2/(j - i + 1)$. That is, $\boldsymbol{E}[X_{ij}] = 2/(j - i + 1)$.

Clearly, $X = \sum_{i,j} X_{ij}$. Hence:

$$\begin{aligned}
\boldsymbol{E}[X] &= \sum_{i,j:i<j} \boldsymbol{E}[X_{ij}] = \sum_{i,j:i<j} \frac{2}{j - i + 1} \\
&= 2\sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \frac{1}{j - i + 1} \\
&= 2\sum_{i=1}^{n-1} O(\log(n - i + 1)) \\
&= 2\sum_{i=1}^{n-1} O(\log n) = O(n \log n).
\end{aligned}$$

Analysis of Quick Sort

The above analysis used the following fact:

$$1 + 1/2 + 1/3 + 1/4 + \ldots + 1/n = O(\log n).$$

The left-hand side is called the harmonic series.