## ENGG1410F Tutorial Introduction to Page Ranks

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Google made its first debut with page ranks, which represent a technique for ranking the webpages on the Internet by importance. Today we will give a short introduction to this technique. Interestingly, at its core, the technique requires computing just an eigenvector.

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Let us model the Internet as a graph. Each webpage is represented as a node. Given two nodes  $v_1, v_2 \in V$ , there is an edge from  $v_1$  to  $v_2$  if the webpage  $v_1$  has a hyperlink to the webpage  $v_2$ .



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- 1. Let  $\boldsymbol{u}$  be a random webpage in the Internet.
- 2. With probability  $\alpha$ :
  - 2.1 If there is at least one out-going link on u2.2 Click on a random hyperlink in u2.3 Set u to the new webpage that opens up.
  - 2.4 Repeat from Step 2.

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- 2.4 Repeat from Step 2.
- 3. With probability  $1 \alpha$ :
  - 3.1 Set u to a random webpage in the Internet.
  - 3.2 Repeat from Step 2.

The value of  $\alpha$  is often set to 0.85 in practice.



For example, suppose we are at  $v_3$ . Conceptually this is what we do:

- Toss a coin that heads with probability  $\alpha$ .
- If the coin comes up heads, jump to  $v_2$  or  $v_4$  with equal chance.
- If the coin comes up tails, jump to  $v_1$ ,  $v_2$ , ...,  $v_5$  with equal chance.

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But this is a wrong question, because the process is random, such that there won't be a deterministic answer. The correct question to ask is:

If the user keeps surfing like this, what is the probability that s/he will land on  $v_1$  as the 10000000000-th page?

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The probability is the page rank of  $v_1$ . Of course, the same question can also be asked about any other page.

If the user keeps surfing like this, what is the probability that s/he will land on  $v_1$  as the 1000000000-th page?

Why the number 10000000000? Interestingly, the theory of random walks (which we will not get into today) tells us that the probability remains the same as long as a sufficiently large number of steps have been performed! In other words, it won't matter if you replace 10000000000 with, say, 100000000001!

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The rationale behind page ranks is this:

A page v is more "important", i.e., having a higher page rank, if a random surfer has a larger chance landing on v after a sufficiently large number of steps.



The page ranks of  $v_1, ..., v_5$  are 0.1716, 0.1666, 0.3214, 0.1666, and 0.1737, respectively. Note that the sum of all the page ranks is 1.

Remaining question: How to calculate them?

Let *n* be the number of nodes. Define  $M = [m_{ij}]$  as an  $n \times n$  matrix where  $m_{ij}$  is the probability of moving from node  $v_j$  to node  $v_i$ .



For example, when  $\alpha = 0.85$ ,  $m_{23} = 0.455$ . Why? See next.

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**Recall**: Suppose we are at  $v_3$ . Conceptually this is what we do:

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So the probability to go from  $v_3$  to  $v_2$  is:

$$\alpha/2 + (1-\alpha)/5$$

which is 0.455 for  $\alpha = 0.85$ .



You can verify:

$$\boldsymbol{M} = \begin{bmatrix} 0.03 & 0.455 & 0.03 & 0.455 & 0.03 \\ 0.03 & 0.03 & 0.455 & 0.03 & 0.03 \\ 0.455 & 0.455 & 0.03 & 0.03 & 0.88 \\ 0.03 & 0.03 & 0.455 & 0.03 & 0.03 \\ 0.455 & 0.03 & 0.03 & 0.455 & 0.03 \end{bmatrix}$$

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Theory of random walks tells us some important facts:

- **M** must have an eigenvalue 1.
- The page ranks make an eigenvector under the eigenvalue 1!

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In our example:

| <b>M</b> =          | 0.03   | 0.455    | 0.03     | 0.455    | 0.03          |       |
|---------------------|--|----------|----------|----------|---------------|-------|
|                     | 0.03   | 0.03     | 0.455    | 0.03     | 0.03          |       |
|                     | 0.455  | 0.455    | 0.03     | 0.03     | 0.88          |       |
|                     | 0.03   | 0.03     | 0.455    | 0.03     | 0.03          |       |
|                     | 0.455  | 0.03     | 0.03     | 0.455    | 0.03          |       |
| You can verify that | 0.1716<br>0.1666<br>0.3214<br>0.1666<br>0.1737 | is indee | d an eig | genvecto | r of <b>M</b> | under |
| eigenvalue 1.       |  |          |          |          |               |       |

We now have an algorithm to compute the page ranks:

- Obtain *M*.
- Obtain an arbitrary eigenvector p of M under the eigenvalue 1.
- Scale p into cp with a proper real number c so that all components of cp add up to 1.

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• *cp* now stores the page ranks of all vertices.