Theory of Computation Complexity

Tutorial 3 LIU yang

Outline

- Data streams
- Counting Distinct Elements
- Sketching

Data Streams

- A data stream is a massive sequence of data.
 - Too large to be stored.

• Examples:

- World Wide Web
- Internet traffic logs
- Financial transactions
- Database transactions

Technique:

- Stream algorithms
- Sketching
- **–** ..

Data Stream Model

- Items in the stream are presented sequentially
- Single pass over the data: i₁,i₂,...,i_n
 - n is known.
- Bounded storage
 - Typically sublinear, n^a or log n.
- Processing time per element must be fast
 - Allow randomness.

E.g. Counting Distinct Elements

- Stream elements: numbers from {1,...,m}
- Goal: Count the number of distinct elements(DE) in the stream
 - Given a T > 0, output YES, if DE>(1+ ϵ)T; No if DE < (1- ϵ)T, with high probability.

Vector Interpretation

Stream: 8 2 1 9 1 9 2 4 4 9 4 2 5 4 2 5 8 5 2 5

- Initially, x=0
- Insertion of i is interpreted as

$$x_i = x_i + 1$$

Want to estimate DE(x)

^{*}This slide is from lecture notes of Prof. Piotr Indyk for the course Sketching, Streaming and Sub-linear Space Algorithms

Estimating DE(x)

Vector X:

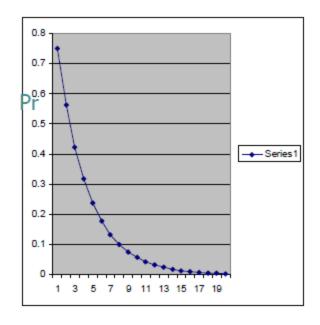


1 2 3 4 5 6 7 8 9

Set S: + + ·

(T=4)

- Choose a random set S of coordinates
 - For each i, we have Pr[i∈S]=1/T
- Maintain Sum_s(x) = Σ_{ies} x_i
- Estimation algorithm A:
 - YES, if Sum_s(x)>0
 - NO, if Sum_S(x)=0
- Analysis:
 - $Pr=Pr[Sum_s(x)=0] = (1-1/T)^{DE}$
 - For T large enough: (1-1/T)DE ≈e-DE/T
 - Using calculus, for ε small enough:
 - If DE> (1+ε)T, then Pr ≈ e^{-(1+ε)} < 1/e ε/3
 - if DE< (1- ϵ)T, then Pr $\approx e^{-(1-\epsilon)} > 1/e + \epsilon/3$



DF

Sketching

- Just like public-coin SMP model
 - Alice and Bob send messages to Referee
 - Referee announces the output
- Streams algorithms can be used for sketching
 - E.g., counting distinct elements
- A simple glimpse
 - Equality test in public-con SMP

Summery

Data Streams

Counting Distinct Elements

Sketching

Thanks