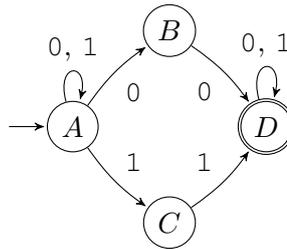


### Week 3 Tutorial Session

1. (a) Write down a regular expression for the following NFA. For this problem, you do not have to go through the procedure described in class.
- (b) Convert the following NFA into a DFA.



2. Let  $L$  be any language. We say that two strings  $x$  and  $y$  are *indistinguishable by  $L$*  if for every string  $z$ , we have  $xz \in L$  if and only if  $yz \in L$ .
  - (a) For concreteness, consider  $L_1 = \{x \in \{0,1\}^* \mid \text{the number of 1's in } x \text{ is divisible by 3}\}$ . Prove that 1 and 1111 are indistinguishable by  $L_1$ .
  - (b) Continuing with (a), which strings are indistinguishable from the string 1 by  $L_1$ ? The set of all such strings is the *equivalence class* of the string 1 and will be denoted by  $[1]$ .
  - (c) Find a string  $s$  not in  $[1]$ . What is the equivalence class of  $s$ ? (We will denote this equivalence class by  $[s]$ )
  - (d) Can you find another string  $t$  not in  $[1]$  or  $[s]$ ? What is the equivalence class of  $t$ ?
  - (e) Can you find yet another string  $u$  not in these equivalence classes?
  - (f) Design a DFA for the language  $L_1$ . How are states in your DFA related to the equivalence classes?
3. (a) Write down the definition of regular expressions over an alphabet  $\Sigma$ .
- (b) Given a string  $w$ , define  $w^R$  as the string  $w$  in reverse order. That is, if  $w = w_1w_2 \dots w_n$ , then  $w^R = w_nw_{n-1} \dots w_1$ . For example, if  $w = \text{live}$ , then  $w^R = \text{evil}$ . Given a language  $L$ , define its reversal  $L^R$  as the set of strings in  $L$  in reverse. More precisely,  $L^R = \{w^R \mid w \in L\}$ . For example, if  $L = \{\text{live, raw, level}\}$ , then  $L^R = \{\text{evil, war, level}\}$ .  
If  $L$  is a regular language, prove that  $L^R$  is also regular.