NFA to DFA conversion and regular expressions

CSCI 3130 Formal Languages and Automata Theory

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DFAs and NFAs are equally powerful

NFA can do everything a DFA can do

How about the other way?

Every NFA is equivalent to some DFA for the same language
Given an NFA, figure out

1. the initial active states
2. how the set of active states changes upon reading an input symbol
NFA → DFA example

Initial active states (before reading any input)?
NFA → DFA example

Initial active states (before reading any input)?

How does the set of active states change?
NFA $\rightarrow$ DFA example

Initial active states (before reading any input)?

partial DFA: $\{q_0, q_1, q_2\}$

How does the set of active states change?
NFA → DFA example

NFA:

Initial active states (before reading any input)?

partial DFA:

How does the set of active states change?
NFA → DFA example

Initial active states (before reading any input)?

How does the set of active states change?
NFA $\rightarrow$ DFA summary

Every DFA state corresponds to a subset of NFA states.
A DFA state is accepting if it contains an accepting NFA state.
Regular expressions
Regular expressions

Powerful string matching feature in advanced editors (e.g. Vim, Emacs) and modern programming languages (e.g. PERL, Python)

PERL regex examples:

\texttt{colou?r} matches “color”/“colour”

\texttt{[A-Za-z]*ing} matches any word ending in “ing”

We will learn to parse complicated regex \textit{recursively}
by building up from simpler ones

Also construct the language matched by the expression \textit{recursively}

Will focus on regular expressions in \textit{formal language theory}
(notations differ from PERL/Python/POSIX regex)
String concatenation

\[ s = \text{abb} \]
\[ t = \text{bab} \]
\[ st = \text{abbbab} \]
\[ ts = \text{bababb} \]
\[ ss = \text{abbabb} \]
\[ sst = \text{abbbabbbab} \]

\[ s = x_1 \ldots x_n, \quad t = y_1 \ldots y_m \]
\[
\Downarrow
\]
\[ st = x_1 \ldots x_n y_1 \ldots y_m \]
Operations on languages

- **Concatenation** of languages $L_1$ and $L_2$
  \[ L_1 L_2 = \{ st \mid s \in L_1, t \in L_2 \} \]

- $n$-th **power** of language $L$
  \[ L^n = \{ s_1 s_2 \ldots s_n \mid s_1, s_2, \ldots, s_n \in L \} \]

- **Union** of $L_1$ and $L_2$
  \[ L_1 \cup L_2 = \{ s \mid s \in L_1 \text{ or } s \in L_2 \} \]
Example

$L_1 = \{0, 01\}$

$L_2 = \{\varepsilon, 1, 11, 111, \ldots \}$

$L_1 L_2 = \{0, 01, 011, 0111, \ldots \} \cup \{01, 011, 0111, 01111, \ldots \}$

$= \{0, 01, 011, 0111, \ldots \}$

0 followed by any number of 1s

$L_1^2 = \{00, 001, 010, 0101\}$

$L_2^2 = L_2$

$L_2^n = L_2$ for any $n \geq 1$

$L_1 \cup L_2 = \{0, 01, \varepsilon, 1, 11, 111, \ldots \}$
The **star** of $L$ are contains strings made up of zero or more chunks from $L$

$$L^* = L^0 \cup L^1 \cup L^2 \cup \ldots$$

Example: $L_1 = \{0, 01\}$ and $L_2 = \{\varepsilon, 1, 11, 111, \ldots\}$

What is $L_1^*$? $L_2^*$?
Example

$L_1 = \{0, 01\}$

$L_1^0 = \{\varepsilon\}$
$L_1^1 = \{0, 01\}$
$L_1^2 = \{00, 001, 010, 0101\}$
$L_1^3 = \{000, 0001, 0010, 00101, 0100, 01001, 01010, 010101\}$

Which of the following are in $L_1^*$?

- 00100001
- 00110001
- 10010001

Yes
No
\[ L_1 = \{0, 01\} \]

\[ L_1^0 = \{\varepsilon\} \]
\[ L_1^1 = \{0, 01\} \]
\[ L_1^2 = \{00, 001, 010, 0101\} \]
\[ L_1^3 = \{000, 0001, 0010, 00101, 0100, 01001, 01010, 010101\} \]

Which of the following are in \( L_1^* \)?

- 00100001 
  Yes

- 00110001 
  No

- 10010001 
  No
Example

$L_1 = \{0, 01\}$

$L_0^1 = \{\varepsilon\}$

$L_1^1 = \{0, 01\}$

$L_2^1 = \{00, 001, 010, 0101\}$

$L_3^1 = \{000, 0001, 0010, 00101, 0100, 01001, 01010, 010101\}$

Which of the following are in $L_1^*$?

- 00100001: Yes
- 00110001: No
- 10010001: No

$L_1^*$ contains all strings such that any 1 is preceded by a 0.
Example

$L_2 = \{ \varepsilon, 1, 11, 111, \ldots \}$

any number of 1s

$L_2^0 = \{ \varepsilon \}$
$L_2^1 = L_2$
$L_2^2 = L_2$
$L_2^n = L_2 \quad (n \geq 1)$
Example

\[ L_2 = \{ \varepsilon, 1, 11, 111, \ldots \} \]

any number of 1s

\[ L_2^0 = \{ \varepsilon \} \]
\[ L_2^1 = L_2 \]
\[ L_2^2 = L_2 \]
\[ L_2^n = L_2 \quad (n \geq 1) \]

\[ L_2^* = L_2^0 \cup L_2^1 \cup L_2^2 \cup \ldots \]
\[ = \{ \varepsilon \} \cup L_2 \cup L_2 \cup \ldots \]
\[ = L_2 \]

\[ L_2^* = L_2 \]
Combining languages

We can construct languages by starting with simple ones, like \{0\} and \{1\}, and combining them

\{0\}(\{0\} \cup \{1\})^* \Rightarrow 0(0 + 1)^*

all strings that start with 0
Combining languages

We can construct languages by starting with simple ones, like \( \{0\} \) and \( \{1\} \), and combining them

\[
\{0\}(\{0\} \cup \{1\})^* \quad \Rightarrow \quad 0(0 + 1)^*
\]

\(0\) followed by any number of 1s, or

\[
(\{0\}\{1\}^*) \cup (\{1\}\{0\}^*) \quad \Rightarrow \quad 01^* + 10^*
\]

0 followed by any number of 1s, or 1 followed by any number of 0s
Combining languages

We can construct languages by starting with simple ones, like \{0\} and \{1\}, and combining them

\{0\}(\{0\} \cup \{1\})^* \quad \Rightarrow \quad 0(0 + 1)^* \\
all strings that start with 0

(\{0\}\{1\}^*) \cup (\{1\}\{0\}^*) \quad \Rightarrow \quad 01^* + 10^* \\
0 followed by any number of 1s, or 1 followed by any number of 0s

\(0(0 + 1)^*\) and \(01^* + 10^*\) are regular expressions

Blueprints for combining simpler languages into complex ones
A language $L$ over $\Sigma$ is **regular** if it is one of the following

- $L = \emptyset$ or $\{\varepsilon\}$
- $L = \{x\}$ where $x$ is a symbol in $\Sigma$
  - If $\Sigma = \{0, 1\}$, then $\{0\}$ and $\{1\}$ are both regular over $\Sigma$
- if $L_1$ and $L_2$ are both regular, so are $L_1 \cup L_2$, $L_1 L_2$ and $L_1^*$
A regular expression over $\Sigma$ is an expression formed by the following rules

- The symbols $\emptyset$ and $\varepsilon$ are regular expressions
- Every symbol in $\Sigma$ is a regular expression
  - If $\Sigma = \{0, 1\}$, then 0 and 1 are both regular expressions over $\Sigma$
- If $R$ and $S$ are regular expressions, so are $R + S$, $RS$ and $R^*$

Examples:

<table>
<thead>
<tr>
<th>$\emptyset$</th>
<th>$\varepsilon$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$0(0 + 1)^*$</td>
<td>$1^*(\varepsilon + 0)$</td>
</tr>
<tr>
<td>$01^* + 10^*$</td>
<td>$(0 + 1)^<em>01(0 + 1)^</em>$</td>
</tr>
</tbody>
</table>

A language is regular if it is represented by a regular expression
Understanding regular expressions

\[ \Sigma = \{0, 1\} \]

\[ 01^* = 0(1)^* \text{ represents } \{0, 01, 011, 0111, \ldots \} \]

0 followed by any number of 1s

\[ 01^* \text{ is not } (01)^* \]
0 + 1 yields \{0, 1\} strings of length 1

\((0 + 1)^*\) yields \{\varepsilon, 0, 1, 00, 01, 10, 11, \ldots\\} any string

\((0 + 1)^*010\) any string that ends in 010

\((0 + 1)^*01(0 + 1)^*\) any string containing 01
What language does the following represent?

$$(((0 + 1)(0 + 1))^* + ((0 + 1)(0 + 1)(0 + 1))^*)$$
What language does the following represent?

\[ (((0 + 1)(0 + 1))^{\ast} + ((0 + 1)(0 + 1)(0 + 1))^{\ast})^{\ast} \]

\[ (((0 + 1)(0 + 1))^{\ast} \quad (0 + 1)(0 + 1)(0 + 1))^{\ast} \]
What language does the following represent?

\[ ((0 + 1)(0 + 1))^* + ((0 + 1)(0 + 1)(0 + 1))^* \]

- \[((0 + 1)(0 + 1))^*\]
- \[((0 + 1)(0 + 1)(0 + 1))^*\]

- \((0 + 1)(0 + 1)\)
- \((0 + 1)(0 + 1)(0 + 1)\)
Understanding regular expressions

What language does the following represent?

\(((0 + 1)(0 + 1))^* + ((0 + 1)(0 + 1)(0 + 1))^*\)

\(((0 + 1)(0 + 1))^*\)

\((0 + 1)(0 + 1)\)
strings of length 2

\((0 + 1)(0 + 1)(0 + 1)\)
strings of length 3
Understanding regular expressions

What language does the following represent?

$$(((0 + 1)(0 + 1))^{*} + ((0 + 1)(0 + 1)(0 + 1))^{*}$$

$$((0 + 1)(0 + 1))^{*}$$
strings of even length

$$((0 + 1)(0 + 1)(0 + 1))^{*}$$
strings whose length is a multiple of 3

$$(0 + 1)(0 + 1)$$
strings of length 2

$$(0 + 1)(0 + 1)(0 + 1)$$
strings of length 3
What language does the following represent?

$$(((0 + 1)(0 + 1)) + (((0 + 1)(0 + 1)(0 + 1))$$

strings whose length is **even or a multiple of 3**

= strings of length 0, 2, 3, 4, 6, 8, 9, 10, 12, . . .

$$(((0 + 1)(0 + 1))$$

strings of **even length**

$$((0 + 1)(0 + 1)(0 + 1))$$

strings whose length is a **multiple of 3**

$$((0 + 1)(0 + 1)$$

strings of length 2

$$((0 + 1)(0 + 1)(0 + 1))$$

strings of length 3
What language does the following represent?

$$((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))^*$$
What language does the following represent?

\(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))\)^* 

\((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1)\)
Understanding regular expressions

What language does the following represent?

$$((0+1)(0+1) + (0+1)(0+1)(0+1))^*$$

$$(0+1)(0+1) + (0+1)(0+1)(0+1)$$

$$(0+1)(0+1)$$

$$(0+1)(0+1)(0+1)$$
What language does the following represent?

$$(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))*)$$

$$((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))$$

$$((0 + 1)(0 + 1)$$

strings of length 2

$$((0 + 1)(0 + 1)(0 + 1)$$

strings of length 3
What language does the following represent?

$$(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))^*$$

$$((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))$$

strings of length 2 or 3

$$(0 + 1)(0 + 1)$$

strings of length 2

$$(0 + 1)(0 + 1)(0 + 1)$$

strings of length 3
Understanding regular expressions

What language does the following represent?

\(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))\)^{*}

strings that can be **broken into blocks**, where each block has **length 2 or 3**

\((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1)\)

**strings of length 2 or 3**

\((0 + 1)(0 + 1)\)

**strings of length 2**

\((0 + 1)(0 + 1)(0 + 1)\)

**strings of length 3**
Understanding regular expressions

What language does the following represent?

\(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))\)^* 

strings that can be broken into blocks, where each block has length 2 or 3

Which are in the language?

ε, 1, 01, 011, 00110, 011010110
What language does the following represent?

\(((0 + 1)(0 + 1) + (0 + 1)(0 + 1)(0 + 1))\)^* 

strings that can be **broken into blocks**, where each block has **length 2** or **3**

Which are in the language?

<table>
<thead>
<tr>
<th></th>
<th>1</th>
<th>01</th>
<th>011</th>
<th>00110</th>
<th>011010110</th>
</tr>
</thead>
<tbody>
<tr>
<td>ε</td>
<td>✔</td>
<td>✗</td>
<td>✔</td>
<td>✔</td>
<td>✔</td>
</tr>
</tbody>
</table>

The regular expression represents all strings except 0 and 1
What language does the following represent?

\[(1 + 01 + 001)^* (\varepsilon + 0 + 00)\]
What language does the following represent?

ends in at most two 0s

\[(1 + 01 + 001)^* \ (\varepsilon + 0 + 00)\]
Understanding regular expressions

What language does the following represent?

\[(1 + 01 + 001)^* (\varepsilon + 0 + 00)\]

- at most two 0s between two consecutive 1s
- ends in at most two 0s
- Never three consecutive 0s

The regular expression represents strings not containing 000

Examples:

\[\varepsilon \quad 00 \quad 0110010110 \quad 0010010\]
Write a regular expression for all strings with two consecutive 0s

\((\text{anything})00\text{anything} \)
Write a regular expression for all strings with two consecutive 0s

(anything)00(anything)

(0 + 1)*00(0 + 1)*