CENG 3420 Computer Organization & Design

Lecture 02: ISA Introduction

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(Textbook: Chapters 1.3 & 2.1)

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Organization – First Glance



Components

- processor (datapath, control)
- input (mouse, keyboard)
- output (display, printer)
- memory (cache, main memory, disk drive, CD/DVD)
- network

Our primary focus: the processor (datapath and control) and its interaction with memory systems

- Implemented using tens/hundreds of millions of transistors
- Impossible to understand by looking at each transistor
- We need abstraction!

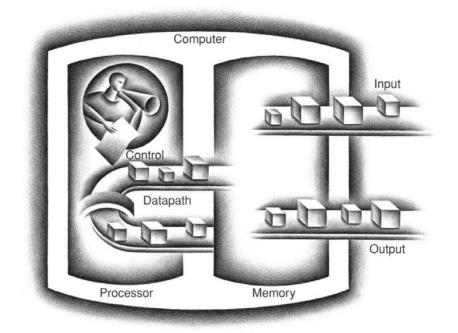


Which one is not a component of a computer?

- A: processor
- B: I/O
- C: memory
- D: fan

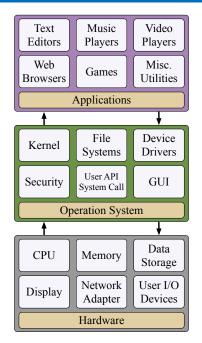
Major Components of a Computer





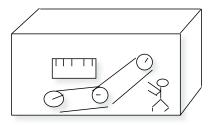
Computer System





Machine Organization

- Capabilities and performance characteristics of the principal Functional Units (FUs). (e.g., register file, ALU, multiplexors, memories, ...)
- The ways those FUs are interconnected (e.g., buses)
- Logic and means by which information flow between FUs is controlled
- The machine's Instruction Set Architecture (ISA)
- Register Transfer Level (RTL) machine description







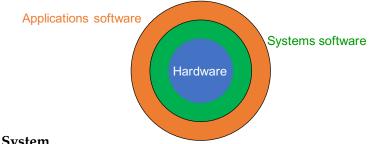
Control needs to have circuitry to

- Decide which is the next instruction and input it from memory
- Decode the instruction
- Issue signals that control the way information flows between datapath components
- Control what operations the datapath's functional units perform

Datapath needs to have circuitry to

- Execute instructions functional units (e.g., adder) and storage locations (e.g., register file)
- Interconnect the functional units so that the instructions can be executed as required
- Load data from and store data to memory





Operating System

- Supervising program that interfaces the user's program with the hardware (e.g., Linux, iOS, Windows)
- Handles basic input and output operations
- Allocates storage and memory
- Provides for protected sharing among multiple applications

Compiler

• Translate programs written in a high-level language (e.g., C, Java) into instructions that the hardware can execute



Which one is not a task of an operating system?

- A: output images to the screen
- B: create files
- C: translate C program to assembly language
- D: switch the processes that run on the processor

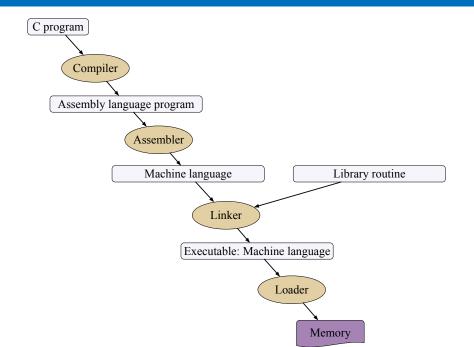


- Allow the programmer to think in a more natural language and for their intended use (Fortran for scientific computation, Cobol for business programming, Lisp for symbol manipulation, Java for web programming, ...)
- Improve programmer productivity more understandable code that is easier to debug and validate
- Improve program maintainability
- Allow programs to be independent of the computer on which they are developed (compilers and assemblers can translate high-level language programs to the binary instructions of any machine)
- Emergence of optimizing compilers that produce very efficient assembly code optimized for the target machine

As a result, very little programming is done today at the assembler level

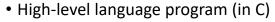
Traditional Compilation Flow





Below the Program





```
swap (int v[], int k)
(int temp;
    temp = v[k];
    v[k] = v[k+1];
    v[k+1] = temp;
)
```



Assembly language program

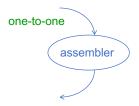
swap: sll \$2, \$5, 2

dd	\$2,	\$4,	\$2
W	\$15,	0(\$	2)
W	\$16,	4(\$	2)
W	\$16,	0(\$	2)
W	\$15,	4(\$	2)
r	\$31		

• Machine (object) code

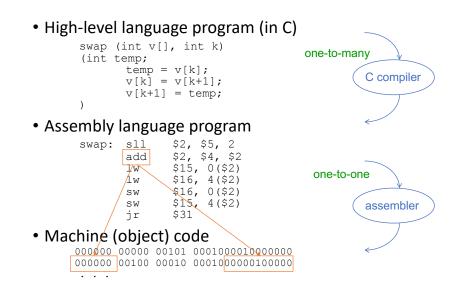
s s i

000000 00000 00101 000100001000000 000000 00100 00010 0001000000100000



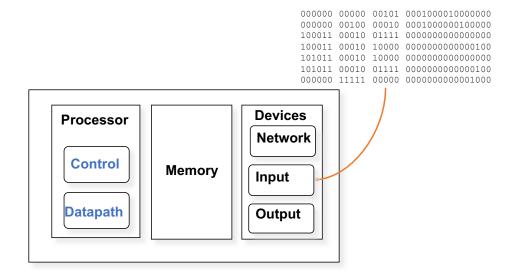
Below the Program



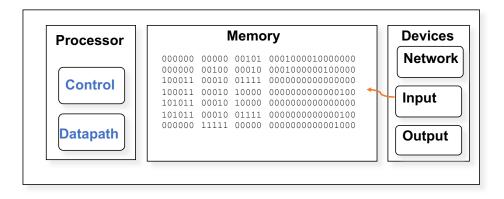


Max # of operations?

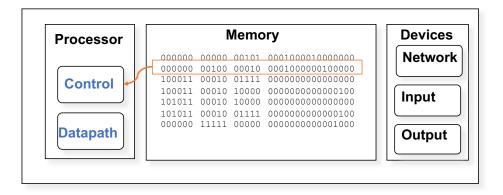






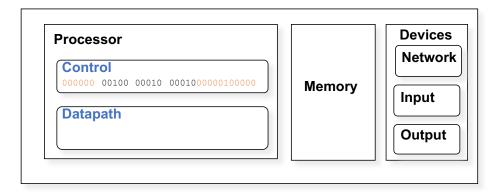






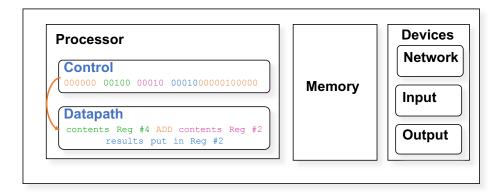
Processor fetches an instruction from memory





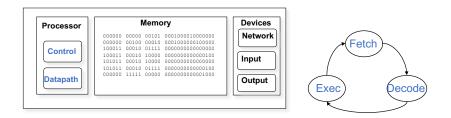
• Control decodes the instruction to determine what to execute



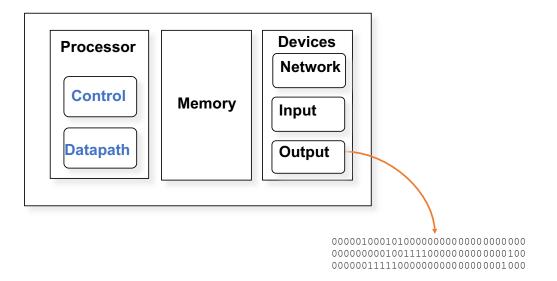


- Control decodes the instruction to determine what to execute
- Datapath executes the instruction as directed by control





- Processor fetches the next instruction from memory
- How does it know which location in memory to fetch from next?

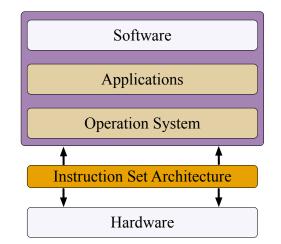




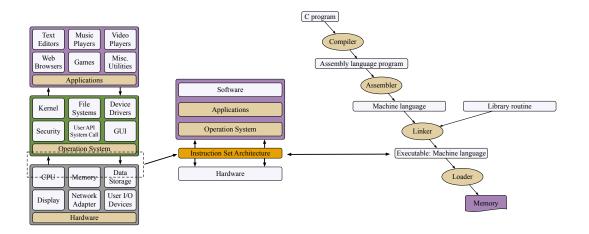








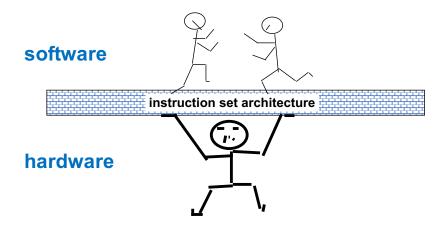






Instruction Set Architecture (ISA)

The interface description separating the software and hardware





- ISA, or simply architecture the abstract interface between the hardware and the lowest level software that includes all the information necessary to write a machine language program, including instructions, registers, memory access, I/O, ...
- Enables implementations of varying cost and performance to run identical software
- The combination of the basic instruction set (the ISA) and the operating system interface is called the application binary interface (ABI)
- ABI: The user portion of the instruction set plus the operating system interfaces used by application programmers. Defines a standard for binary portability across computers.

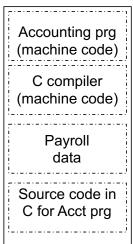
Two Key Principles of Machine Design



- 1) Instructions are represented as numbers and, as such, are indistinguishable from data
- Programs are stored in alterable memory (that can be read or written to) just like data Memory

Stored-Program Concept

- Programs can be shipped as files of binary numbers binary compatibility
- Computers can inherit ready-made software provided they are compatible with an existing ISA leads industry to align around a small number of ISAs



The language of the machine

• Want an ISA that makes it easy to build the hardware and the compiler while maximizing performance and minimizing cost

Our target: the **RISC-V** ISA

- similar to other ISAs developed since the 1980's
- RISC-V is originated from MIPS, the latter of which is used by Broadcom, Cisco, NEC, Nintendo, Sony, ...

Design Goals

Maximize performance, minimize cost, reduce design time (time-to-market), minimize memory space (embedded systems), minimize power consumption (mobile systems)







Complex Instruction Set Computer (CISC)

Lots of instructions of variable size, very memory optimal, typically less registers.

• Intel x86

Reduced Instruction Set Computer (RISC)

Instructions, all of a fixed size, more registers, optimized for speed. Usually called a "Load/Store" architecture.

• RISC-V, LC-3b, MIPS, ARM, Sun SPARC, HP PA-RISC, IBM PowerPC ...

- Used in many embedded systems
- E.g., Nintendo-64, Playstation 1, Playstation 2









RISC-V

- An open standard instruction set architecture (ISA)
- A clean break from the earlier MIPS-inspired designs
- Modular ISA organization
- Open standards, numerous proprietary and open-source cores
- Managed by RISC-V Foundation



Instruction Categories

- Load and Store instructions
- Bitwise instructions
- Arithmetic instructions
- Control transfer instructions
- Pseudo instructions

4 Base Instruction Formats: all 32 bits wide

31	25 2	24	20 19	15	14 1	2 11	76	0
funct	7	rs2	r	s1	funct3	rd	opcode	R-type
iı	nm[11:0]		r	s1	funct3	rd	opcode	I-type
imm[11	:5]	rs2	r	s1	funct3	imm[4:0] opcode	S-type
		imm[31:1	[2]			rd	opcode	U-type
L								



Which one is not a part of ISA?

- A: load/store instructions
- B: arithmetic operations
- C: printf function
- D: control tranfer instructions



Which one is not a part of ISA?

- A: load/store instructions
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- D: control tranfer instructions

Ans: C



Register Names	ABI Names	Description	
x0	zero	Hard-wired zero	
x1	ra	Return address	
x2	sp	Stack pointer	
x3	gp	Global pointer	
x4	tp	Thread pointer	
x5	t0	Temporary / Alternate link register	
x6-7	t1 - t2	Temporary register	
x8	s0 / fp	Saved register / Frame pointer	
x9	s1	Saved register	
x10-11	a0-a1	Function argument / Return value registers	
x12-17	a2-a7	Function argument registers	
x18-27	s2-s11	Saved registers	
x28-31	t3-t6	Temporary registers	



Which register cannot we store an operand when we call "max(a, b)"?

- A: x0 (zero)
- B: x3 (gp)
- C: x9 (s1)
- D: x11 (a1)



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- A: x0 (zero)
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- D: x11 (a1)

Ans: A