

Problem 1

Design a TM M to accept the language $L = \{0^n 1^n \mid n \geq 1\}$.

Solution

Initially, the tape of M contains $0^n 1^n$ followed by an infinity of blanks. Repeatedly, M replaces the leftmost 0 by X , moves right to the leftmost 1, replacing it by Y , moves left to find the rightmost X , then moves one cell right to the leftmost 0 and repeats the cycle. If, however, when searching for a 1, M finds a blank instead, then M halts without accepting. If, after changing a 1 to a Y , M finds no more 0's, then M checks that no more 1's remain, accepting if there are none.

Let $Q = \{q_0, q_1, q_2, q_3, q_4\}$, $\Sigma = \{0, 1\}$, $\Gamma = \{0, 1, X, Y, B\}$, and $F = \{q_4\}$. Informally, each state represents a statement or a group of statements in a program. State q_0 is entered initially and also immediately prior to each replacement of a leftmost 0 by an X . State q_1 is used to search right, skipping over 0's and Y 's until it finds the leftmost 1. If M finds a 1 it changes it to Y , entering state q_2 . State q_2 searches left for an X and enters state q_0 upon finding it, moving right, to the leftmost 0, as it changes its state. As M searches right in state q_1 , if a B or X is encountered before a 1, then the input is rejected; either there are too many 0's or the input is not in $0^* 1^*$.

State q_0 has another role. If, after state q_2 finds the rightmost X , there is a Y immediately to its right, then the 0's are exhausted. From q_0 , scanning Y , state q_3 is entered to scan over Y 's and check that no 1's remain. If the Y 's are followed by a B , state q_4 is entered and acceptance occurs; otherwise the string is rejected. The function is shown below.

| | 0 | 1 | X | Y | B |
|-------------------|---------------|---------------|---------------|---------------|---------------|
| $\rightarrow q_0$ | (q_1, X, R) | — | — | (q_3, Y, R) | — |
| q_1 | $(q_1, 0, R)$ | (q_2, Y, L) | — | (q_1, Y, R) | — |
| q_2 | $(q_2, 0, L)$ | — | (q_0, X, R) | (q_2, Y, L) | — |
| q_3 | — | — | — | (q_3, Y, R) | (q_4, B, R) |
| $* q_4$ | — | — | — | — | — |

Problem 2

Design Turing machines to recognize $\{ww^R \mid w \text{ is in } (0+1)^*\}$

Solution

| | 0 | 1 | X | B |
|-------------------|---------------|---------------|---------------|---------------|
| $\rightarrow q_0$ | (q_1, X, R) | (q_2, X, R) | (q_6, X, R) | (q_6, B, R) |
| q_1 | $(q_1, 0, R)$ | $(q_1, 1, R)$ | (q_4, X, L) | (q_4, B, L) |
| q_2 | $(q_2, 0, R)$ | $(q_2, 1, R)$ | (q_5, X, L) | (q_5, B, L) |
| q_3 | $(q_3, 0, L)$ | $(q_3, 1, L)$ | (q_0, X, R) | — |
| q_4 | (q_3, X, L) | — | — | — |
| q_5 | — | (q_3, X, L) | — | — |
| * q_6 | — | — | — | — |

The state q_0 goes right on the tape and find the first one that is not X , say a , replace it by X . Then goes right to find the first X or B , if its left symbol is a , replace it by X , otherwise, reject.