## Exercises for CSCI5010

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**Problem 1.** Let P be a set of points in  $\mathbb{R}^d$ . Consider the following algorithm for finding an approximation of diam(P) (i.e., the diameter of P).

## algorithm

- 1. pick an arbitrary point  $p \in P$
- 2. identify a point  $q \in P$  maximizing dist(p,q) (the Euclidean distance between p and q)
- 3. return dist(p,q)

Prove: The algorithm returns a  $\frac{1}{2}$ -approximate answer, i.e.,  $\frac{1}{2} \cdot \operatorname{diam}(P) \leq \operatorname{dist}(p,q) \leq \operatorname{diam}(P)$ .

**Problem 2.** Let *P* be a set of points in  $\mathbb{R}^d$ . Describe how to find in  $O(n \log n + n/\epsilon^d)$  time a value  $\Delta$  satisfying diam $(P) \leq \Delta \leq (1 + \epsilon) \cdot \text{diam}(P)$  for any  $0 < \epsilon < 1$ .

**Problem 3\*.** Let P be a set of 2D points. Let  $p, q \in P$  be two points such that dist(p,q) = diam(P). Prove: p and q must be vertices of the convex hull of P.

Hint: Use the fact that if a point p is not a vertex of the convex hull, then every line  $\ell$  passing p has the property that there are points of P falling on both sides of  $\ell$ .

**Problem 4\*.** Let G be a convex polygon. Define p and q as the two vertices of G maximizing dist(p,q). Prove: G has an edge e such that

- one of p and q is a vertex of e;
- if p is a vertex of e, then q is a vertex of G having the maximum distance to the line  $\ell$  passing e.



Hint: First find two parallel lines passing p and q, respectively, and enclosing G in between. Then, rotate these lines.

Remark: Problems 3 and 4 suggest that the (precise) diameter of a 2D point set P can be found in  $O(n \log n)$  time.

**Problem 5 (reading exercise).** Prove the correctness of the algorithm discussed in the lecture for computing a *t*-spanner graph for a set of points in  $\mathbb{R}^d$ .

Hint: If you do not want to think, read Pg 128-129 of the lecture notes. However, an inductive argument similar to the one for solving Problem 4 of Exercise List 6 suffices.

**Problem 6 (reading exercise).** Prove the correctness of the algorithm discussed in the lecture for computing a  $(1 + \epsilon)$ -approximate EMST for a set of points in  $\mathbb{R}^d$ .

Hint: If you do not want to think, read Pg 129-130 of the lecture notes.